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[54] **COMPUTER SYSTEM HAVING SYSTEM FEATURE EXTENSION SOFTWARE CONTAINING A SELF-DESCRIBING FEATURE TABLE FOR ACCESSING I/O DEVICES ACCORDING TO MACHINE-INDEPENDENT FORMAT**

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[51] Int. Cl.⁵ **G06F 9/445; G06F 9/44; G06F 9/00**

[52] U.S. Cl. **395/500; 395/800; 395/700; 364/280; 364/280.2; 364/280.9; 364/DIG. 1; 364/975.2; 364/976**

[58] Field of Search **395/800, 700, 500**

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Primary Examiner—Thomas C. Lee

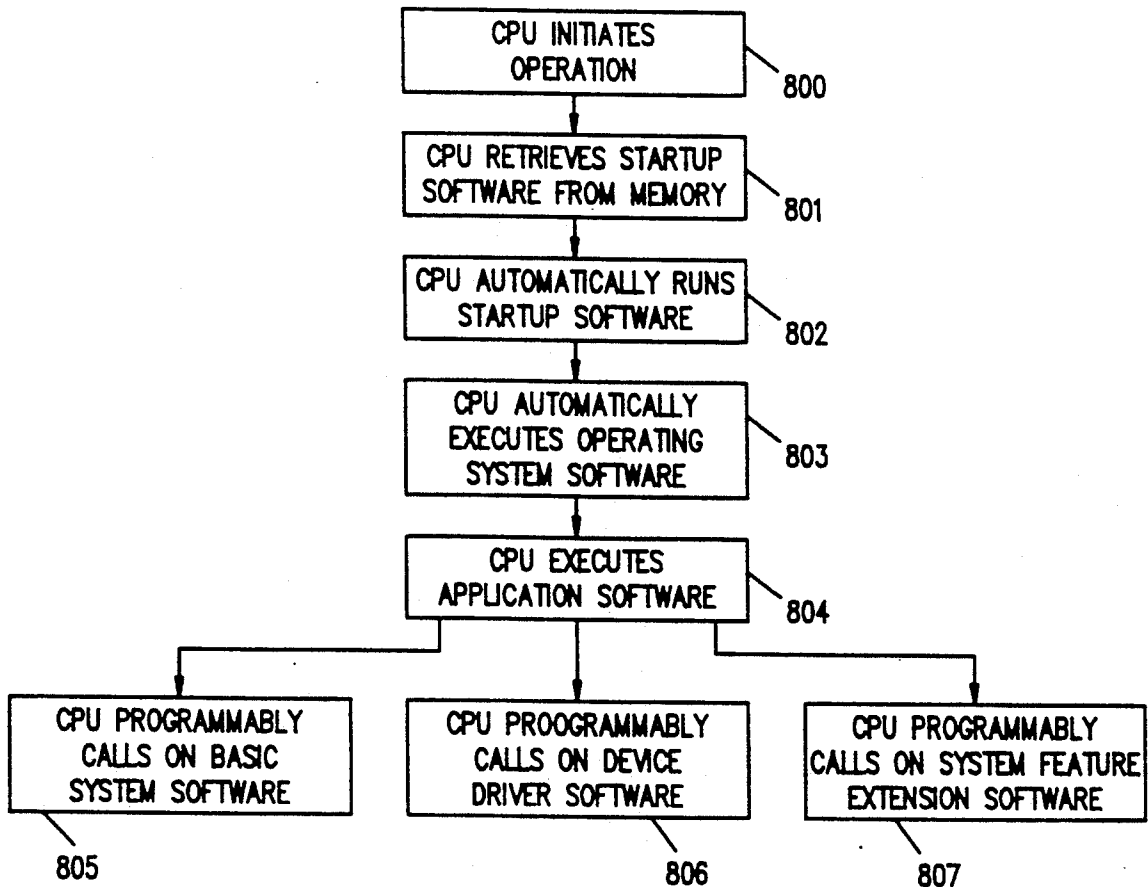
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[57] **ABSTRACT**

A computer system with self-describing feature table, accessible by device drivers. Thus a simple process can access these feature tables to fully customize the device drivers at installation, or at boot; or the device driver can branch on the data in the feature table. Thus, a new degree of flexibility is achieved without degrading performance.

46 Claims, 6 Drawing Sheets



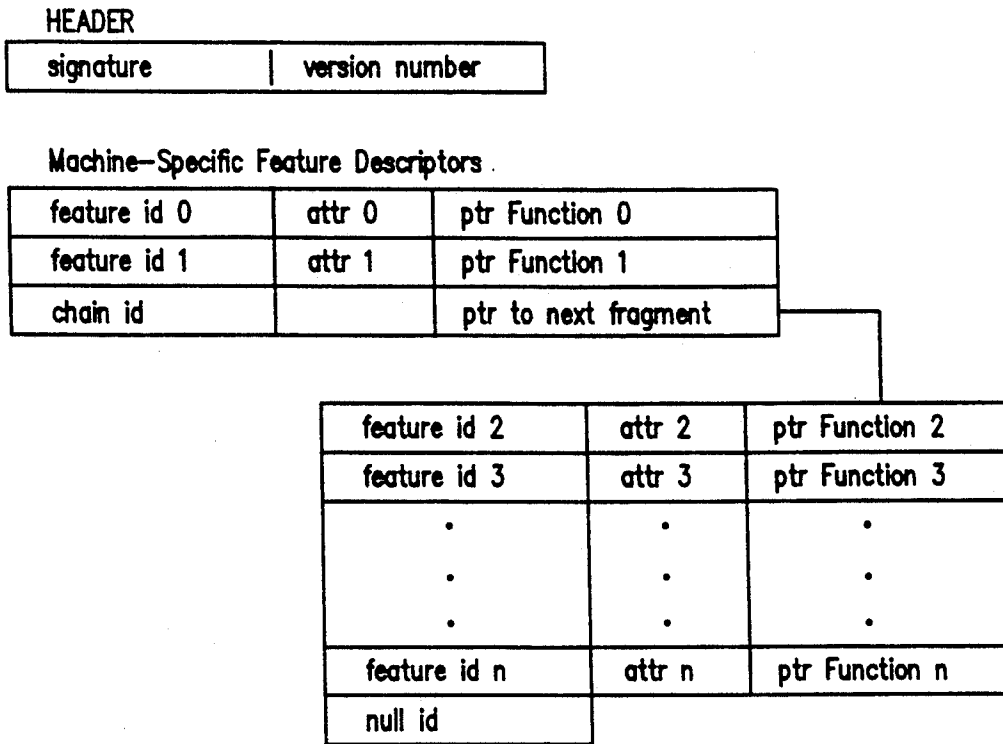
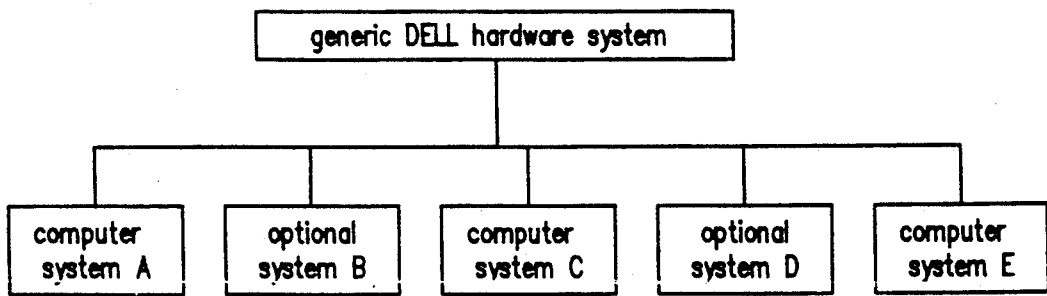


Figure 1



Hardware system class hierarchy.
The generic class object will contain a table of hardware features which each system subclass can inherit or modify.

Figure 2

MEMBER	LENGTH	DESCRIPTION
feature id	2 bytes	Selector used to identify system features
attributes	2 bytes	Feature characteristic indicators where: Bit 0 = XBIOS ptr is a real mode entry Bit 1 = XBIOS ptr is a protect mode entry Bimodal = real & protect bits set Data = real & protect bits zero Bit 2 = Reserved Bit 3 = XBIOS ptr format 0 = linear 1 = seg:offset Bit 4 = optional keystroke trigger Bit 5 = optional appendix Bits 6-15 = Reserved
XBIOS ptr	4 bytes	References a feature routine or data
keystroke	2 bytes	Optional feature keystroke trigger where: High byte = keystroke trigger shift state Bit 0 = Ctrl shift state Bit 1 = Alt shift state Bits 2-7 = Reserved Low byte = scan code
appendix	variable	Optional data append to entry where: Length = 2-byte field containing the size of the subsequent variable block Block = variable length appended data

Figure 3

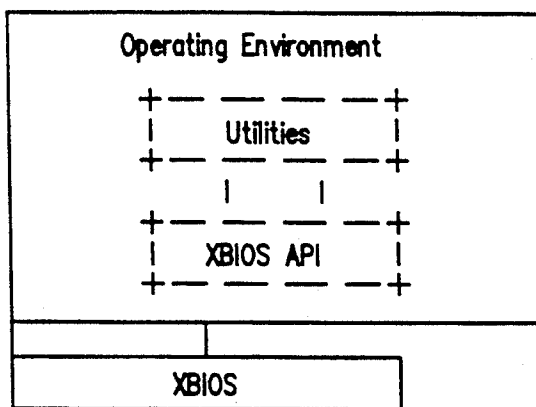
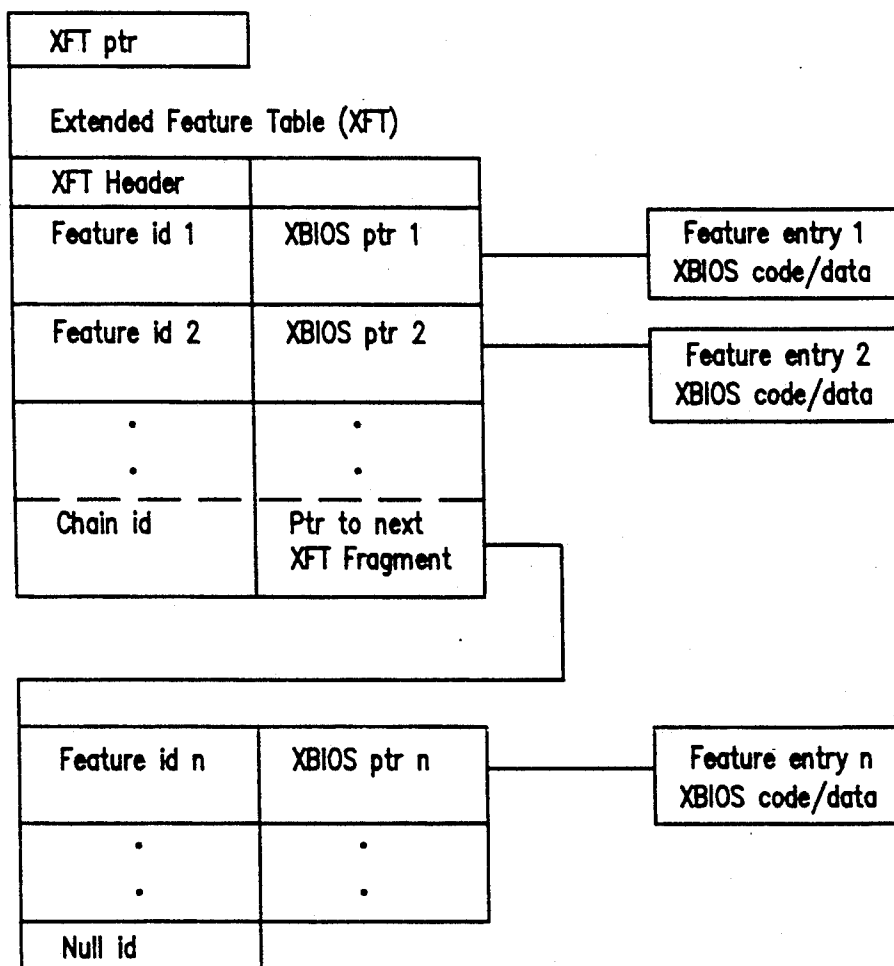


Figure 4



* F000:ED00 (XBIOS Fixed address - start of OEM reserved area)

Figure 5

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;* XBIOS FEATURE ROUTINE TEMPLATE *;

xbinput      struc                ;input map
    xi0      dw    ?
    xi1      dd    ?
    xi2      db    ?
xbinput      ends

xboutput     struc                ;output map
    xo0      dw    ?
xboutput     ends

xbstkframe   struc                ;stack frame map
    xin      db    (type xbinput) dup (?) ;input area
    xout     db    (type xboutput) dup (?) ;output area
    xsub     dw    ?                ;subfunction id
xbstkframe   ends

XbiosProc    proc    far
    push    ds                ;save ds register . . .
    push    cs                ;set ds register . . .
    pop     ds                ; addressability
    .
    mov     cx, es:[bx] .xin.xi0 ;get 1st input parm
    mov     es:[bx].xout.xo0, cx ;move cx to output
    .
    mov     si, es:[bx] .xsub    ;get subfunction id . . .
    shl    si, 1                ; determine address . . .
    call   subfunctbl [si]      ;Call the subfunction.
    .
    mov     ax, XBS_STAT        ;move status to AX.
    pop     ds                ;restore ds register
    ret
XbiosProc    endp

```

Figure 6

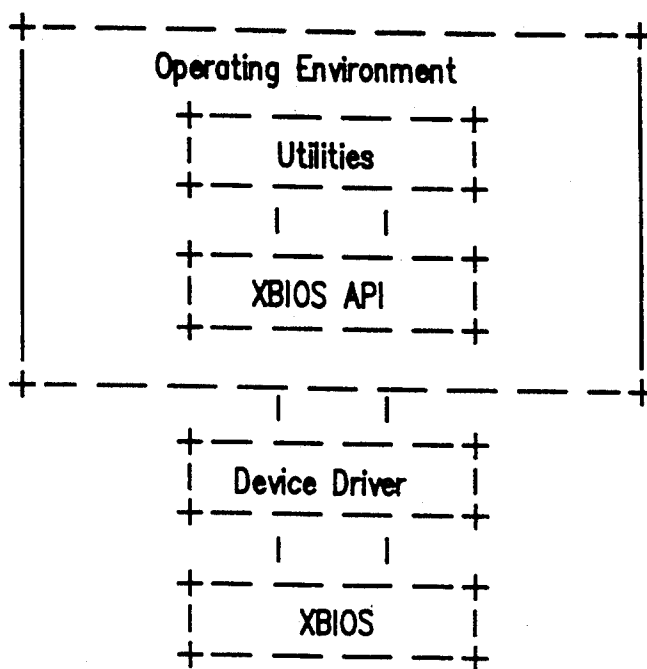


Figure 7

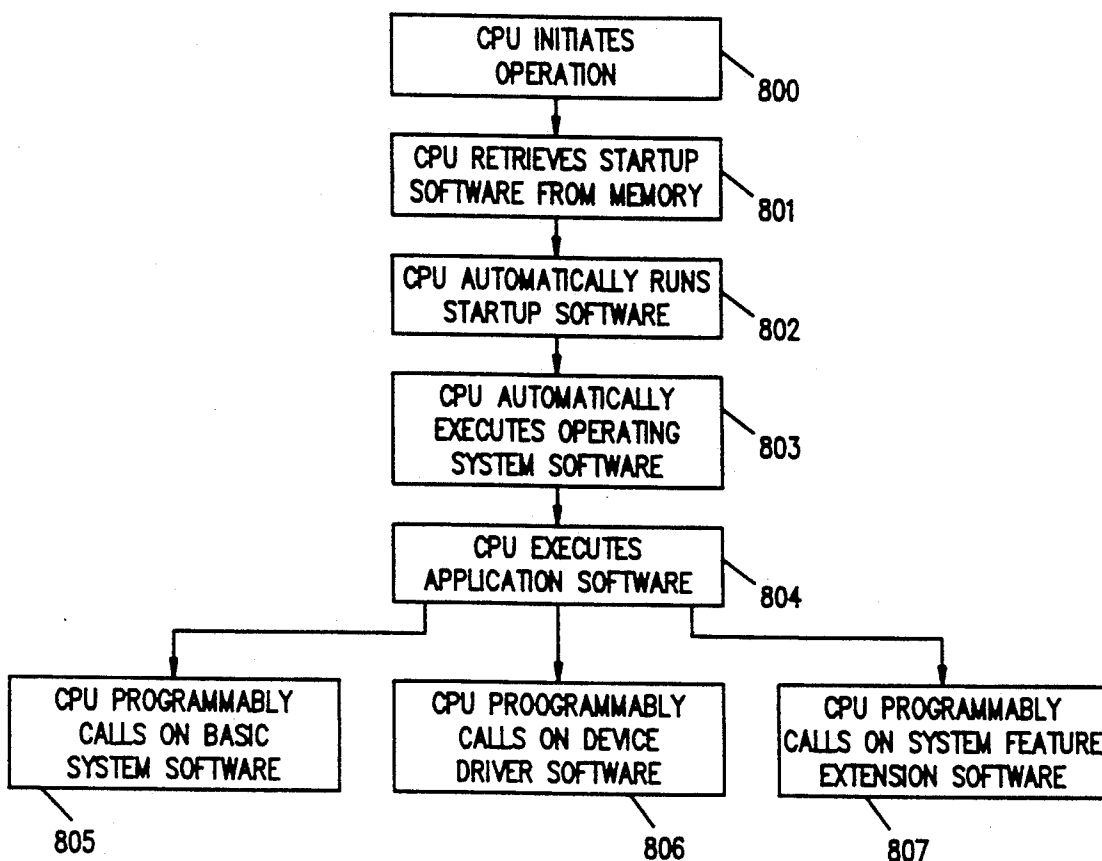


Figure 8

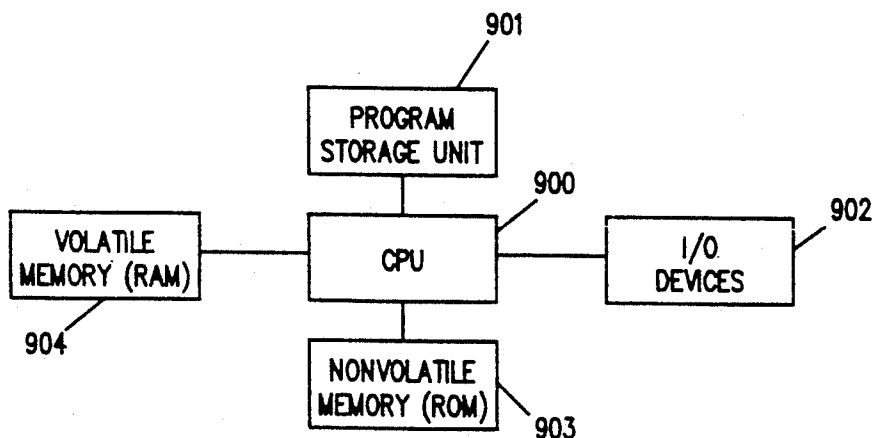


Figure 9

**COMPUTER SYSTEM HAVING SYSTEM
FEATURE EXTENSION SOFTWARE
CONTAINING A SELF-DESCRIBING FEATURE
TABLE FOR ACCESSING I/O DEVICES
ACCORDING TO MACHINE-INDEPENDENT
FORMAT**

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**BACKGROUND AND SUMMARY OF THE
INVENTION**

The present invention relates to computer systems, and particularly to single-user or few-user small systems.

How Application Programs Interact with Hardware

One of the most basic needs in computer architecture is making it easier for a variety of software programs to interact correctly and efficiently with a variety of hardware configurations. Much of the development in computer architecture can be seen as a steady progression of techniques for addressing this need.

Note the emphasis on efficiency in the foregoing statement. Even where existing standards can assure compatibility, the search for greater speed or expanded functionality will frequently draw programmers to circumvent the software standards. A good example of this countercurrent appeared in the early days of graphics development on the IBM PC: the BIOS provided a standard interface to video driver operations, but software developers discovered that they could vastly improve performance by making calls directly to the video driver hardware. Thus adherence to the standard architecture was not enough to assure computer designers that their customers would be able to run popular IBM-compatible software, such as Flight Simulator™, on their supposedly IBM-compatible machines. Thus, there is a continuing tension between compatibility and efficiency.

When any particular piece of hardware is examined in isolation, it can usually be best described in terms of electrical relationships. For example, a memory specification may state that, within a certain range of delay after certain voltages appear on certain lines, certain other lines will be driven to a corresponding state (which is dependent on the data previously stored in the memory). The specification for an input device, such as a keyboard, may state that, when certain voltages appear on certain lines, a particular input operation may be considered to have occurred. The specification for

an output device, such as a video card, may state that, when certain voltages appear on certain lines within a certain timing relationship and protocol, each pixel within a certain defined display device will be driven to an optical state which corresponds to a certain portion of the protocol. However, a commercial application program will be written in a programming language (e.g. in assembly language or in C) which is somewhat machine-independent. There is a great difference between these two levels of description; but this gap must be bridged in order to economically develop application software which can run on a wide range of machines.

Several layers of software and firmware structure are used to mediate between application programs and the underlying hardware. To better show the context of the invention, these layers will be described below in greater detail.

Hardware Variability

Computer hardware configurations are inherently diverse. The complexity of any modern computer system is high enough that even a very detailed standard architecture (such as the "AT" architecture which was introduced with IBM's 80286-based machines) will not prevent variations from occurring. Whenever designers independently work within a standard, they are likely to find ways to make improvements. As such variations occur, some of them will be seen to be significant. Thus the future will often find that any standard contained significant "gray" areas.

This is true not only in motherboard design, but also in I/O devices. For example, two display drivers which both conform to the VGA standard may nevertheless differ in timing, to an extent which may be significant to some software applications. Moreover, there will always be users with needs for specialized input or output devices.

Even within the very restricted world of "PC" architectures (where all machines must conform to numerous constraints of the "standard" architectures), hardware variability continues to be a problem. The range of hardware variability in (for example) computers which can run UNIX is far larger.

This hardware variability is not merely accidental, but will continue: users are eager to take advantage of new developments, and the pace of innovation is generally far too rapid to permit stabilization of standardized hardware configurations.

Layers of Software and Firmware Structure

In order to mediate between application programs and the underlying hardware, several layers of software and firmware structure are used. To better show the context of the invention, these layers will be described below in greater detail.

Startup Software (POST, Bootstrap, etc.)

When power is first applied to a computer, the various hardware elements (chips and subsystems) will each have their own internal procedures (reset procedures) to regain a stable and known state. However, at some point (if the hardware is intact), these reset procedures will have ended, and at this point the CPU performs various important overhead tasks. These include, for example, surveying the system configuration, performing sanity checks on system hardware, issuing diagnos-

tic signals (such as sounding beeps through a speaker or turning on LEDs), and permitting the user to branch into an NVRAM configuration program under software control. This phase of operation is generally referred to as "POST" (Power-On-Self-Test). After POST, a "bootstrap" program is run, to permit the CPU to begin execution of other software. For robustness, the POST and bootstrap software is normally stored in a read-only memory. The bootstrap program launches the CPU on execution of the primary operating system software; Depending on how the system has been set up, the boot software may direct program execution into DOS, Unix, PS/2, a DOS variant, or another operating system. This is normally automatic and predetermined, but is manually user-selectable in some systems. However, the choice of operating system is not particularly relevant to the inventions described in the present application, the primary operating system can then be used by the user to launch an application program, either manually or automatically.

Basic Input/Output System Software (BIOS)

In many types of modern personal computers (and in all "IBM-compatible" personal computers), a key part of the system software is a "basic input/output system" (BIOS) program. See generally, e.g., the P. Norton, *THE PETER NORTON PROGRAMMER'S GUIDE TO THE IBM PC* (1985), which is hereby incorporated by reference. The BIOS program contains frequently-used routines for interfacing to key peripherals. The term "peripheral" or "peripheral component" normally refers to those components of a computer system which are not on the motherboard, i.e. which must be addressed through a system bus or over a defined port. However, the usage of this term is somewhat variable; sometimes it is used to refer to any I/O device, or only to refer to components which are optional add-ons. For interrupt handling, and so forth. Thus, the BIOS software provides some degree of machine-independence. However, in PC-class computers, this independence is not fully exploited by the available commercial software. Many programs bypass the BIOS software, and directly access the underlying hardware addresses or devices. See generally Glass, "The IBM PC BIOS," *BYTE*, April 1989, pp. 303ff.

For system robustness, the BIOS software is normally packaged in a read-only-memory. However, in 1991 IBM introduced a PS/2 system in which the BIOS is at least partially stored on disk. In fact, it is normally packaged together with the startup software mentioned above. Packaging the BIOS, POST and boot routines in ROM makes a very robust firmware system. Short of hardware damage, it is very difficult for a user to distort the system to the point where it will not start up and run (if the operating system software is present). However, this system also provides a considerable degree of flexibility. As the operating system up (after the POST and boot routines), the user can remap address pointers to revector BIOS calls away from the standard BIOS routines, if desired. (It is also common for users to map out the entire BIOS contents into fast RAM, for greater speed). Thus, nowadays the term "BIOS" is often used, somewhat more broadly, to refer to this whole collection of basic system routines. However, in the present application references to "BIOS" will normally refer to the BIOS in its narrower sense, i.e. to the collection of I/O handling routines (and associated routines) which

can be called on by the operating system or by the application software.

Customized BIOS and BIOS Extensions

The BIOS in IBM-compatible computers is accessed by interrupts, but the vectors for those interrupts can be diverted to other addresses (by overwriting an address pointer in system RAM). This capability significantly expands the flexibility of the BIOS, and programmers use it very frequently.

However, while the capability to divert BIOS vectors is useful, it is not sufficient to address many needs. Changes to the interrupt-handling vectors will not affect other portions of the BIOS. Computer designers have found it highly desirable to prepare (or obtain) customized BIOS routines to fully exploit the advantage of their systems. For example, such customized BIOS routines are commonly necessary in very-low-power portable systems, to implement power-saving features which maximize battery lifetime. BIOS customization has increasingly been recognized as an important element in rapidly developing a reliable advanced system. See generally Scheier, "Phoenix counters competitors with diversified BIOS offerings," *PC Week*, vol. 4 no. 38 (Sep. 22, 1987) at 135f; Guterman, "CompuAdd adopts new ROM BIOS for clones," *PC Week Vol. 5 no. 28* (Jul. 11, 1988) at 6; both of which are hereby incorporated by reference.

One function often provided by BIOS customization is "hot-key" access to a setup menu, or to low-level system hardware features (e.g. monitor brightness adjustment). Such capability is very useful to system designers, but normally it has had to be realized in a machine-dependent way (so that large chunks of BIOS have had to be rewritten every time a change was made).

Another problem with prior hot-key add-ons is that, if the BIOS interrupt vector for key-handling was diverted, the hot-key capability could be lost. Since many applications do divert the keyboard interrupt (INT9), no critical functionality could be made dependent on such a hot-key operation.

Operating System Software

The application software will normally interface to an operating system (such as DOS, DOS+Windows, OS/2, UNIX of various flavors, or UNIX plus X-windows). The operating system is a background software program. Some operating systems run continuously, or at least start up at regular intervals, even while an application program is running; other operating systems merely provide capabilities which can be called on by the application software. which provides an application programming interface (API) for use by the application software. Thus, the programmers writing application software can write their software to fit the API, rather than having to find out and fit the peculiarities of each particular machine. See e.g., Quedens, "Windows virtual machine," *PC Tech Journal* vol. 5, no. 10 p. 90, 92-3, 95, 97, 99-100, 102 (Oct. 1987), which is hereby incorporated by reference.

Graphical User Interface (GUI) Operating System Add-Ons

Some operating systems have been enhanced by the addition of overlaid supplemental operating systems. For example, Windows is a supplement to DOS, and X is a supplement for UNIX. The use of such hybrids does

not greatly affect the foregoing considerations, except that it makes the compatibility issues even more difficult: the designer of a DOS machine must expect that customers will be running some DOS programs, and some Windows programs, on the same machine.

Device Driver Software

A device driver is a lower level of operating system software. Typically a device driver interfaces to the actual peripheral hardware components, and provides routines which application software can use to access the hardware components. Thus, the application software can simply make a call to an installed software subroutine, instead of having to find the specifications of each peripheral device and branch accordingly, whenever a peripheral I/O operation is needed. This permits application software to ignore the detailed specifications of peripheral hardware.

Normally device driver software must contain a description of each target hardware platform. Thus, the software must be revised repeatedly, for reasons which are beyond the control of the companies making peripherals.

In personal computers, installable device drivers were first introduced in DOS 2.0. The role of device drivers has since been expanded, in subsequently introduced operating systems.

In particular, OS/2 provided expanded support for device drivers, including a library of "DevHlp" routines which can be called by device drivers. See generally Duncan, "An examination of the DevHlp API (writing OS-2 bimodal device drivers)," 3 Microsoft Systems Journal no.2 (March 1988) at 39ff; Schmitt, "Designing drivers for OS/2: I," PC Tech Journal vol.5, no. 12, p. 164 (1987); and Schmitt, "Designing drivers for OS/2: II," PC Tech Journal vol.6, no.2 p. 136-155 (Feb. 1988), all of which are hereby incorporated by reference.

System Configuration Tables

Some computer systems have previously used a feature table, stored in nonvolatile memory, to describe various characteristics of the machine. The IBM AT BIOS uses such a feature table (stored in battery-backed CMOS memory). This feature table, in expanded form, has also been used in the IBM PS/2 systems and has been utilized in the system BIOS of all IBM AT- and PS/1-compatible personal computers. This table is in the form of a bit map where each bit refers to specific hardware implementations employed by the designers of the machine. A pointer to this table may be obtained through executing a software interrupt. More specifically, executing interrupt 15h with AH=COH will return a pointer to the table in ES:BX. However, this feature table is restricted to merely listing certain hardware features in the machine, such as the number of DMA controllers, and does not provide an interface to these features. Furthermore, the elements of the list are fixed.

In the Phoenix Technologies BIOS, there are specified entry points at hard-coded addresses which will perform certain machine-specific functions. These are few in number, must be at fixed addresses, do not support protected mode applications and it is not possible to easily see which features are supported by which machines except by restricting to the common subset.

Under the EISA standard, an EISA Configuration memory is used to store a limited feature table of EISA

peripherals on the EISA bus. See generally Glass, "Inside EISA," BYTE magazine, November 1989, at 417ff, which is hereby incorporated by reference.

Application Software

From a system designer's point of view, the application software is (subject only to the minimal constraints of the architectural standards) wholly unpredictable. Many clever people are constantly looking for new ways to exploit the standard architecture, and many innovations continually result. Thus, hardware architects must expect that the application software will not only be unpredictable, but will be as unpredictable as possible. Common applications include spreadsheets, word processors and publishing programs, databases, games, project managers and a wide variety of others; but inevitably users will also run customized applications, and new types of applications.

Utility Programs and Hardware

In recent years, many personal computer manufacturers have expanded their product lines. This has dramatically increased the difficulty of supporting an entire product line in terms of the standard software products that a manufacturer may choose to include or sell with its computers.

Examples are diagnostic programs, operating system software and utility software. It is increasingly necessary to provide a means for such software to identify the individual machines and their unique features, without having to be rewritten each time a new product is introduced.

Furthermore, it may be difficult or undesirable to implement even similar features in exactly the same way, since each design has different constraints in terms of cost and each will incorporate the knowledge gained by building the previous product. The problem gets worse as a product line ages. It is desirable to continue to support older products with newer versions of software, and it is also desirable for older versions of software to run unmodified on newer platforms. One solution to this problem is to write the software to the common subset of functions supported by all platforms. However, this does not allow the manufacturer to differentiate his product from the competition. Consequently, it is desirable for each individual machine to have the capability to identify its own unique feature set to such software, while at the same time providing the individualized means for carrying out those functions.

Innovative Computer System with Self-Describing Extensions to BIOS

The present invention provides a personal computer architecture with an additional layer of overhead software (or firmware) structure. This additional layer of software structure is used to provide access to additional low-level hardware-specific features in a manner which is independent of the operating system. In the present application, these additional low-level hardware-specific features are referred to as "extended features."

Extended Features

An "extended feature", in the presently preferred embodiment, is normally a system level routine used to service hardware components or to obtain system information unique to Dell hardware systems. The detailed disclosure below lists some of the numerous extended

features which have been implemented to date. However, of course, other functions can be provided as well.

The disclosed self-describing system software extension provides a lower level of software-hardware interface "cushioning," which device drivers can call on. Thus, the self-describing system software extension can also be exploited to permit device drivers to be more hardware-independent.

The self-describing system software extension is particularly advantageous in its application to an evolving product line within the same overall standard.

Self-Describing Feature Table

The disclosed innovations provide methods by which a computer with some quantity of non-volatile storage can present a self-describing interface which also provides a means for carrying out machine-specific functions in a non-specific way.

In the presently preferred embodiment, the feature table and the machine-specific routines are programmed into EPROM devices, at the start of the "OEM reserved" block of addresses in the BIOS memory space. In IBM-compatible computers, the BIOS commonly occupies the 64K or 128K of address space just below the top of the lowest megabyte of the total memory address space.

An important element of the method is a table which contains a signature to identify it with the system software architecture described herein, and a series of entries with the following information:

Feature ID—a unique identifier for each specific machine-specific function.

Attributes—describe the operating environment for proper access to the function. May limit access to real or protected mode or possibly even to specific operating environments.

Service Routine—a pointer to the program code that performs the requested function.

Data Block—Features may also include an optional data block.

The self-describing system software extension of the presently preferred embodiment includes a self-describing feature table, which can track the peculiarities of the actual hardware configuration of each system as configured. The self-describing system software extension, with this feature table, provides low-level translation for hardware peculiarities.

By use of this feature table, the disclosed innovations provide a computer system which can be updated with self-defining extensions to the basic BIOS (which remains in read-only memory). The basic BIOS must be modified to make use of these self-defining extensions; but, once such a modified BIOS has been installed, it does not have to be updated frequently. Instead, the routines in the modified BIOS can make use of the self-defining feature table without further changes to other portions of BIOS. Thus, for example, in one class of alternative embodiments, the feature table is located in NVSRAM, and a ROM holds the basic BIOS and a pointer to the feature table.

Application Programming Interface to Self-Describing Feature Table

One contemplated and advantageous class of embodiments uses a standard API to the feature table to provide increased portability (across applications) of access to the extended features. (Thus, for example, an OS/2 device driver can be written to wrap this API around

calls to the feature table in such a way that any OS/2 software can make feature-table calls through this API.) This provides optimal access to the machine-specific routines across the whole family of computers.) In the sample embodiment described in detail below, this function is not yet included. However, as will be apparent to those skilled in the art, this can readily be implemented in various ways, within the architecture described below, if desired.

Device Drivers in the Innovative System

The disclosed architecture provides access to machine-specific features, with enough information to permit device drivers to be written in a machines-independent way (within the Dell family of computers). Some specific examples of such drivers are given herein, but of course other drivers can also make use of the extended BIOS features as well.

Keyboard Driver

One very advantageous piece which is included in the presently preferred embodiment is the keyboard driver, which permits the user to access extended system functions, without exiting his application, by hitting "hot key" combinations. This is highly advantageous in portable systems, since the user can fine-tune his system's hardware parameters to match changing conditions. Thus, the disclosed system allows the user to send BIOS-level function calls right through an application, without saving the context of the application.

This keyboard driver runs under DOS, so it is still possible for a user to disrupt this driver by remapping the keyboard interrupt (INT9); but at least this driver does permit language customization to be combined with hot-key access to extended-BIOS functions.

An important point is that, even when a user remaps the keyboard interrupt under BIOS, he can still preserve the hot-key calls to the extended BIOS features without knowing what the features are or even which key combinations call them. (This is accomplished, in the presently preferred embodiment, by building a quick-reference table in system RAM during an initialization phase.)

OS/2 Initialization Driver

Another advantageous part of the presently preferred embodiment is an OS/2 initialization driver, which permits easy initialization of hardware-specific functions for OS/2 initialization. This driver, in its presently preferred embodiment, is listed in the appendix.

Common Device Drivers across a Family of Computers

The disclosed innovations have been implemented on a number of different computer systems within the Dell system product line. As of the effective filing date of the present application, these include all Dell computer models shipped after September 1990, plus a few models which were retrofitted.

As discussed below, the disclosed innovations are believed to be advantageous not only as applied to a single computer, but also as applied to a whole family of computers. In the case of the preferred embodiment, suppliers of computer peripherals can be increasingly confident that a device driver which takes advantage of the self-describing system software extension of the presently preferred embodiment will apply to every current Dell computer, and also to future models which the supplier has not yet seen or heard of.

Backward Software Compatibility

A substantial advantage of the disclosed architecture is that additional BIOS-level functions can be readily added into computer system designs, as soon as the innovations occur, without any necessity for radical BIOS changes. The self-describing system software extension of the presently preferred embodiment itself does not degrade BIOS compatibility with prior ISA or EISA machines; and once the self-describing system software extension of the presently preferred embodiment is installed, further extensions to BIOS functionality can readily be achieved.

Additional Background

Two previous proposals for achieving machine-independence will now be discussed, with reference to the present inventions, in order to provide a clearer discussion of how the teachings of the present application differ from these prior teachings.

The "Advanced BIOS" (ABIOS) in PS/2 Architectures

The PS/2 architecture, which IBM introduced in 1987, included a so-called "advanced BIOS" or "ABIOS" There are actually two versions of ABIOS available, since IBM has offered a simplified ABIOS for use on machines other than IBM PS/2s. However, the full functionality of ABIOS is available only on an IBM PS/2. The features of this version of ABIOS are most germane to the background of the present invention. (In addition to a more conventional BIOS, known as the "compatibility BIOS" or "CBIOS"). The user can elect to use the CBIOS instead of the ABIOS if he wishes, for downward compatibility; the CBIOS and ABIOS are not designed to run simultaneously.

The ABIOS is a more high-level software structure than ordinary BIOS, and has many features added to enhance performance in OS/2 (which, unlike DOS, is a multi-threaded operating system). However, ABIOS is so complex and ambitious that very few operating system designs have used it.

The ABIOS must normally be initialized: the initialization process surveys the system configuration, and builds a data structure (the CDA) in System RAM. Process threads (or system software called by process threads) can call on this data structure to get information about the hardware they are running on.

The ABIOS is somewhat analogous to a large-scale machine-specific device driver: a process thread can make calls to the ABIOS by submitting a "request block" into the ABIOS's request-handling queue.

When running on an IBM PS/2 under OS/2, the OS/2+ABIOS combination does make additional DevHlp functions available to device drivers, including provision of a standardized interface which provides some hardware-independence to the device driver software. Thus, the device driver software programs in such a system can include substantially increased functionality. See generally Mizell, "Understanding device drivers in Operating System/2," IBM Systems Journal vol.27, no.2 p.170-84 (1988), which is hereby incorporated by reference.

Within the IBM PS/2 family, the interface to ABIOS is identical, regardless of which IBM PS/2 machine it is running on. Within the IBM PS/2 family, the ABIOS provides a significant degree of machine-independence.

Thus the ABIOS has some, but not all, of the same goals as the presently preferred embodiment disclosed herein.

Note that the self-describing system software extension of the presently preferred embodiment provides a lower-level component of system software than the ABIOS referenced above. The ABIOS is itself a full standalone BIOS, and may be thought of as a fancy device driver. By contrast, the conventional BIOS is interrupt-driven. By contrast, the self-describing system software extension will not work as a standalone BIOS, and does not even work as a device driver: instead, the self-describing system software extension merely provides services to device drivers and to standalone application programs.

The self-describing system software extension provides functions which are not addressed by the ABIOS, and conversely the ABIOS addresses a great deal of functionality which is not addressed by self-describing system software extension. Thus, these two software systems are complementary. In fact, it would readily be possible to prepare a modification of self-describing system software extension for use as an ABIOS extension. The ABIOS also includes "hooks" for extending ABIOS; insofar as known to the present inventors, nobody has ever taken the trouble to implement this in a practical system, but there is no apparent reason why this could not be done if desired.

The disclosed self-describing system software extension also provides particular advantages in system diagnostics, which are not provided by ABIOS.

The "XBIOS" of the Atari ST

In internal documentation, the self-describing system software extension of the presently preferred embodiment has frequently been referred to as the "XBIOS," and reflections of that terminology can be seen in the source code appended to the present application. However, to prevent confusion, it must be noted that the term "XBIOS" has also been used for a component of the operating system software in the Atari ST computer. While this software is believed not to have included any of the innovative concepts claimed herein, the similarity in terminology should be noted. See generally, e.g. Rothman, "Atari ST software development, BYTE magazine, vol. 11 no. 9 (Sep. 1986) at223ff, which is hereby incorporated by reference. The Atari ST is a 68000-based machine. The ST's operating system (called "TOS") has two main parts: The "GEM" (Graphics Environment Manager) is a complete operating system developed by Digital Research, and is meant to support applications that are portable to other machines. Atari's "XBIOS" (extended BIOS) is meant to support ST-specific capabilities not accessible through GEM.

Additional Background Literature

U.S. Pat. No. 4,589,063, which is hereby incorporated by reference, purports to disclose a "method and apparatus for automatic configuration of a computer system . . . wherein one or more system peripheral or I/O devices can be interfaced to the computer system through I/O boards that plug into a system motherboard. Each of the I/O devices includes a controlling device driver module that operates under a program code stored in a read only memory resident on the I/O board and by which the device driver module allows the computer system to communicate with its associated peripheral and I/O devices. Accordingly, a system

user is not required to change the computer operating system kernel to support each new I/O device or system configuration change."

BRIEF DESCRIPTION OF THE DRAWING

The present invention will be described with reference to the accompanying drawings, which show important sample embodiments of the invention and which are incorporated in the specification hereof by reference, wherein:

FIG. 1 diagrams the layout of the extended feature table (XFT) used in the presently preferred embodiment.

FIG. 2 schematically shows an object-oriented paradigm wherein, by viewing every DELL hardware system as a subclass of a generic DELL hardware system class, supporting all extended features across every systems, BIOS programmers can choose to inherit or modify any or all of these features for their particular system.

FIG. 3 shows additional details of the structure of each individual feature entry.

FIG. 4 shows how a pointer is used to allow BIOS programmers to relocate the feature table, and enables uniform access to the table regardless of the Dell system type.

FIG. 5 shows how the feature table interface in the application interface (API) library, in the presently preferred embodiment, locates and executes self-describing system software extension feature routines.

FIG. 6 shows how the system-software-extension API library, in the presently preferred embodiment, communicates with the device driver via the I/O control API supplied by the operating environment.

FIG. 7 shows a typical system-software-extension feature routine organization.

FIG. 8 diagrams a flow chart of an implementation of the preferred embodiment of the invention.

FIG. 9 illustrates a preferred embodiment of the present invention.

DESCRIPTION OF THE PREFERRED EMBODIMENTS

The numerous innovative teachings of the present application will be described with particular reference to the presently preferred embodiment. However, it should be understood that this class of embodiments provides only a few examples of the many advantageous uses of the innovative teachings herein. In general, statements made in the specification of the present application do not necessarily delimit any of the various claimed inventions. Moreover, some statements may apply to some inventive features but not to others.

In particular, the following text frequently references the "XBIOS" system software, which is the presently preferred embodiment of the claimed system extension software. The following text also makes frequent references to Dell™ computers, since the preferred embodiment of the system extension software has been implemented in this line of computers. However, of course, the claimed inventions can readily be adapted to a tremendous range of computers and of software implementations.

Even the specific disclosed embodiment is not inherently limited to Dell computers, but provision for other makers' computers, of comparable architecture, can readily be added.

Preferred Embodiment: Family of Computers with Shared General Architecture

The disclosed innovations have been implemented on a number of different computer systems within the Dell system product line. As of the effective filing date of the present application, these include all Dell computer models shipped after September 1990, plus a few models which were retrofitted. The specific models include at least Dell models 325D, 333D, 433P, 333P, 325P, 425E, 433E, 425TE, 433TE, 433DE, 450SE, 433SE, 420DE, 420SE, 320LT, 320N, and 212N. These models include tower, desktop, laptop, and notebook models; EISA- and ISA-bus systems; systems based on 80486, 80386, and 386SX microprocessors; systems running at 50 MHz, 33 MHz, 25 MHz, 20 MHz, and 12 MHz clock rates; systems with one or two hard disks, or up to 10 disks in a drive array; systems with monochrome, VGA, or high-resolution graphics adapters; and a wide variety of other configuration options. Moreover, the disclosed innovations are currently being made available in all new Dell computer designs. As of the filing date of the present application, every computer which is currently being shipped in the Dell product family includes a version of the XBIOS described below. (Of course, this is not necessarily true of the many older models which are already in service in the field.)

As discussed below, the disclosed innovations are believed to be advantageous not only as applied to a single computer, but also as applied to a whole family of computers. In the case of the preferred embodiment, suppliers of computer peripherals can be increasingly confident that a device driver which takes advantage of the self-describing system software extension of the presently preferred embodiment will apply to every current Dell computer, and also to future models which the supplier has not yet seen or heard of.

A sample source code implementation is set forth in the Appendix to insure the fullest possible compliance with the best mode requirements of U.S. patent law. Although the sample source code does represent the state of this code as of the effective filing date of the present application, it must be noted that this specific example is still under development. It is expected that this source code will later be modified to add functionality, improve performance, and possibly also to remove bugs.

To give a more clear example of the workings and advantages of the disclosed innovative system and method ideas, the following is an excerpt from an OS/2 loader that uses several XBIOS features. This code includes the following functions (in several modules):

```

check for existence of XBIOS table in system BIOS
(EPROM);
Looks for SYSTEM_IDENTIFY routine (find xbios
routine);
Call SYSTEM_IDENTIFY routine in EPROM;
Save the data returned by SYSTEM_IDENTIFY
for later use;
Look for OS/213 INIT routine (in EPROM), and call
it if it exists;
Look for the SMARTVU routine, and, if it exists,
display "OS/2" on the diagnostic display;
Looks for GATEA20 routine, and, if it exists, save
the GATEA20 address (from table) for later use.
```

Part of the source code shown on Appendix Page A-ii checks for the existence of an XBIOS table.

The source code shown on Appendix Pages A-ii and A-iii checks for the existence of a System Identify Routine.

The source code shown on Appendix Pages A-iii and A-iv checks for the existence of a SmartVu routine, and uses it, if it exists, to display "OS/2".

The source code shown on Appendix Pages A-iv and A-v calls the OS/2 initialization routine, if it exists.

The source code shown on Appendix Pages A-v and A-vi checks for the existence of a Gate-A20 routine, and stores its address, if it exists, for use in switching between real and protected modes.

The source code shown on Appendix Pages A-vi and A-vii scans the XBIOS table in EPROM looking for a requested function id.

To give a more clear example of the workings and advantages of the disclosed innovative system and method ideas, the source code shown on Pages A-vii to A-xii is an excerpt from a DOS keyboard driver.

The source code shown on Appendix Page A-xii swaps the INT 9 vectors.

The source code shown on Appendix Pages A-xiii to A-xiv checks for a key chord which would require a call to XBIOS.

The source code shown on Appendix Pages A-xiv to A-xvi actually handles keystrokes as desired.

To give a more clear example of the workings and advantages of the disclosed innovative system and method ideas, the source code shown on Appendix Pages A-xvi to A-xviii is an excerpt from a DOS memory manager (HIMEM.SYS) that uses XBIOS functions. It looks for the GATE A20 routine in XBIOS and saves the address for later use by the operating system.

To give a more clear example of the workings and advantages of the disclosed innovative system and method ideas, the source code shown on Appendix Pages A-xviii to A-xxvi is a sample XBIOS table definitions in Assembler.

To give a more clear example of the workings and advantages of the disclosed innovative system and method ideas, the source code shown on Appendix Pages A-xxvi to A-xxviii is a sample of XBIOS table definitions in C.

To give a more clear example of the workings and advantages of the disclosed innovative system and method ideas, the source code shown on Appendix Pages A-xxviii to A-xxxvii is a sample of XBIOS test code (written in C). This is an 80386-based 33 MHz desktop ISA machine, with a typical configuration of a 200M IDE disk drive, 4M of DRAM, 64K cache SRAM, and an 8 MHz ISA bus. However, as detailed above, the disclosed innovations have been implemented on many other computers too.

System Software Extension—Technical Specification

The presently preferred embodiment provides a family of "IBM-compatible" computers. In this family, the disclosed innovations are applied to augment the ROM BIOS by self-defining ROM BIOS feature extensions in a manner independent of the operating environment and extends the system software support for disparate hardware features in a standardized fashion. This strategy is the culmination of ideas arising from the necessity to make the access procedure to the BIOS uniform, and thereby reduce the number of releases of Dell supported operating environment, to provide support for extended system features across all system software platforms, to accommodate enhanced diagnostic sup-

port, and to obtain standardized access to system services.

Dell computer systems support various ROM BIOS (hereafter referred as "BIOS") extensions that enhance the standard AT architecture providing added value to Dell's customers. Hitherto, a keyboard interface and various DOS utilities have been provided for Dell's customers to access these extended features. However, this access procedure is inadequate for computer systems sold overseas and various operating environments. The keyboard interface is sufficient for DOS only systems sold in the US since these systems, by default, use BIOS to process keystrokes that access extended features. Systems sold outside the US, however, use a memory resident DOS utility that traps the BIOS INT 9 (keyboard handler) routine in order to process requested services. This utility requires frequent modifications and testing whenever new features are added to hardware systems since the supporting code is imbedded within the utility.

Moreover, operating systems and graphical environments such as OS/2, WINDOWS, and UNIX intercept keystrokes and process them in a manner which bypass the BIOS keyboard handler. Thus the extended features are not available to users via the keyboard in these environments.

Some utilities written for DOS require direct BIOS accessibility, and hence cannot execute in operating environments that prevent direct access. In order to support compatibility with current systems that access extended features via the keyboard, operating environments that place restrictions on BIOS accessibility, and the additional requirement to improve system diagnostic support, a standardized access method to BIOS extended features independent of the operating environment and hardware system is provided.

The disclosed architecture permits access to extended BIOS features independently of the operating environment. This is accomplished by letting the features be "self-defining". This refers to the ability of the operating environment to access features in a manner independent of the hardware system. An extended feature is viewed by the operating environment as an abstract hardware device or service processed by BIOS with its characteristics embedded in BIOS. The operating environment's only reference to the extended feature is through an identifier that is subsequently defined by BIOS.

The centerpiece of this architecture is a table embedded in BIOS that contains a list of extended features supported by the hardware and its respective attributes. Through the use of this table, the operating environment and the extended features are totally isolated from one another. This allows the operating environment to access extended features in a consistent manner independent of the BIOS. Conversely, the extended features can be developed independently of the specific operating environment, since they are redefined by the BIOS in each specific machine.

The table improves upon the current software interrupt access method to BIOS, since this is not supported by all operating environments. OS/2, for example, does not allow this. DOS has a different problem: programmers can generally redefine any software interrupt, which means that access to extended feature support could be cut off by applications or even by users.

By using the disclosed innovations, BIOS programmers can modify extended features without affecting

other parts of the architecture. This reduces the need to update vendor source code. Utilities that take advantage of these extended features can be developed independently of the operating environment, and thus will have greater portability.

Object-Oriented Paradigm for Feature Inheritance

By utilizing the embedded BIOS table as a method table, the architecture can be developed around an object-oriented paradigm to access the extended features in BIOS. As shown in FIG. 2, by viewing every Dell hardware system as a subclass of a generic Dell hardware system class, supporting all extended features across every systems, BIOS programmers can choose to inherit or modify any or all of these features for their particular system. (Thus, this is only a two-level hierarchy, which avoids problems with "grandchild inheritance.") Each extended feature is assigned a unique identifier. When utility applications need to access this feature, they can use this unique identifier in a message that is dispatched to an interface. This interface subsequently determines the behavior of the extended feature by matching the message to the identifier within the table.

The Self-describing System Software Extension

The "XBIOS" self-describing system software extension contains hardware specific features that extend the standard BIOS operations. An "extended feature" is either a system level routine or data used to service hardware components or to obtain system information unique to Dell hardware systems. Since XBIOS is an extension to BIOS, access to all the standard BIOS functions will remain intact via the interrupt vector table, but an additional access method is provided for the extended features. To implement this access scheme, in the presently preferred embodiment, a table called the extended feature table ("XFIT") is embedded in BIOS. The XFT is used to match service requests for extended features from the operating environment to identifiers representing the extended features listed in the table. Corresponding attributes associated with each feature identifier determine whether the extended feature is an XBIOS function or a pointer to a block of data. The XF-F permits BIOS programmers to define, modify and support extended features for any system without having support built directly into the operating environment or utilities.

Extended Feature Table ("XFT")

The XFT is a table containing extended system features supported by XBIOS. FIG. 1 diagrams the layout of the XFT. The table consist of a header followed by an array of extended hardware feature entries that can be fragmented throughout XBIOS. The header is composed of a signature used to detect XBIOS support by verifying the existence of the XFT, and the XFT version number.

FIG. 3 shows additional details of the structure of each individual feature entry. Each entry contains at least three fixed members: a feature ID, attribute flags, and a 32-bit XBIOS pointer. Each table may also contain two optional members (as determined by the attribute flag settings): a feature keystroke trigger and an appendix-block of additional data.

The XFT is relocatable, as are all the XBIOS extended features. Anchoring an XFT pointer at the same XBIOS address in all systems provides a standard

method of finding the XFT regardless of the system, and allows BIOS programmers the flexibility to locate the table anywhere in XBIOS.

XFT Header

Signature: At header offset 0, the signature ("DELLXBIOS") is a null terminated string of bytes used to identify the existence of an XFT confirming that XBIOS is supported by the system.

XFT Version Number: At header offset 2, the XFT version is a 2-byte value used to verify the current XFT version.

XFT Feature Entry

Feature ID: At feature offset 0, the feature ID is a 2-byte value serving as a selector to identify system supported features. The XFT is scanned to match a feature request against the XFT feature ID list. If a match is found then the request is processed. No match indicates that the request is unsupported. Two feature ID's are reserved for XFT support operations: Chain ID and NULL ID. The Chain ID (0×FFFF) is used to indicate that the corresponding 32-bit pointer references the next table fragment in XBIOS. The NULL ID (0×0000) is used as the XFT termination entry.

Attributes: At feature offset 2, the attributes is a 2-byte bit field containing various characteristics about the feature entry. The flags describes the type of the corresponding XBIOS pointer, whether the features supports a DOS compatible keystroke trigger, and determines if the feature entry record contains extended information. The attribute flags are described as follows:

real mode: 1 bit. Real mode code pointer flag.

protect mode: 1 bit. Protected mode code pointer flag.

XBIOS ptr format: 1 bit. Denotes linear or seg:offset. keystroke trigger: 1 bit. Optional keystroke trigger. appendix flag: 1 bit. Feature entry has appended data.

To indicate that the corresponding XBIOS pointer references bimodal code, both real mode and protected mode flags are set. If neither the flags are set then it is assumed that the XBIOS pointer references data and an XBIOS address is returned. The pointer format flag indicates the address format. When set, the pointer format is segment:offset; when clear, the pointer format is linear.

XBIOS Pointer: At feature offset 4, the XBIOS pointer is a 32-bit pointer that references either code or data depending upon the corresponding attribute flag settings. If the real mode or protected mode flag is set then the pointer contains an XBIOS feature routine entry point. If the data mode is set then the pointer simply contains an address to a block of data. The XBIOS pointer is either a linear address or an address of segment:offset form. Whenever there is a conflict between the usage of either address format, segment:offset should be used since conversion to linear form will always yield a valid address.

Keystroke: A keystroke trigger is used to maintain compatibility with systems that can access features through the keyboard. Whenever the keystroke trigger attribute flag is set a two byte keystroke field immediately follows the permanent members of the feature entry at offset 8. The keystroke field is a 2-byte field consisting of the keyboard shift state in the high byte and the scan code in low byte. For example, if a keystroke trigger is designated as ctrl-alt-enter, the ctrl and

alt state bits are set in the high byte and low byte value is set to hex 1C (enter key scan code).

Appendix: The appendix contains supplemental information attached to the feature entry. This provides greater control and flexibility to XBIOS feature design. The appendix follows either the permanent members of the feature entry record at offset 8, or the keystroke field at offset 10 (if one is designated by the keystroke trigger attribute flag). The first two bytes of the appendix contain the length of the subsequent data block.

XFT Pointer

The XFT pointer is a 32 bit pointer in segment:offset format anchored at the start of the OEM reserved area—location F000:ED00—in YBIOS. This allows BIOS programmers to relocate the XFT and enables uniform access to the table regardless of the Dell system type (FIG. 4). XFT isolates BIOS programmers and reduces the impact from changes made to XBIOS. This isolation allows BIOS programmers to continue to deploy the development environment that comply best with their needs. XFT requires only the extended features supported by the system and BIOS programmers can add or remove features as desired.

Standard and Generic Interface Configurations

The XBIOS interface is organized into two configurations: standard and generic. In the standard XBIOS interface configuration, which is used in the presently preferred embodiment, direct access to XBIOS is provided by the XBIOS API library that is linked to the utility application. The library contains an XFT interface that locates XBIOS features and executes XBIOS feature routines (FIG. 5). In the generic XBIOS interface configuration, access to XBIOS is accomplished indirectly through a device driver that contains the XFT interface. The XBIOS API library communicates with the device driver via the I/O control API supplied by the operating environment (FIG. 6).

Generic Interface through API

In this alternative version, access to XBIOS is provided via an application programming interface (API) used by utilities. Utilities are applications that interact with users and need control of the system-dependent features. The XBIOS API corresponds to the features provided by XBIOS and is consistent among all operating environments and Dell computer systems. By restricting the access to XBIOS only through the API, portable system dependent utilities can be developed in a machine independent style with a high level language. The utilities thus developed will be portable to other operating environments and to other Dell systems. This also permits the XBIOS interface to be organized into various configurations based upon the strategy that best supports the operating environment and customers needs.

Access to XBIOS Feature Routines

XBIOS feature routines are machine specific functions embedded in the ROM BIOS of each hardware system. The XFT interface executes feature routines indirectly using the XBIOS feature pointer when either the real or protect attribute flag is set. The routines must adhere to a standardize XBIOS function call protocol. This protocol enables the XFT interface to call any XBIOS function in a uniform manner.

Each routine defines a set of input and output variables that are passed to and from the feature routine via a parameter buffer and returns a status value back to the XFT interface. Routines can optionally define subfunctions under a single feature ID which are executed through a subfunction identifier also placed in the parameter buffer. XBIOS internal variables are defined within the routine's code segment (usually in segment F000h), and addressability is obtained by assigning the data segment register (DS) to the code segment register (CS).

Upon entry, the XBIOS routine assigns DS to CS (after saving DS on the stack), and receives a pointer to the parameter buffer in ES:BX. The input variables within the parameter buffer are addressed incrementally from ES:BX followed by the output variables and the optional subfunction identifier. Various status flag are passed to the routine in the AX register that can be used to convey information such as the processor modes (real/protected; USE16/JLJSE32). One alternative which was dropped from the presently preferred embodiment was to use the system XFT interface stack for parameter variables. This would have allowed a "C" language interface to the XBIOS routines. Due to the possible stack addressing discrepancies from the base pointer (BP) in the USE16 and USE32 address modes of the processor, the current model using ES:BX was chosen. However, as 32-bit architecture and operating environments become increasingly standard, it may be advantageous to implement such alternative XBIOS routines to support the "C" language interface.

Upon exit, the XBIOS routine returns successful (zero) or failure (non-zero) status in AX and restores the DS register. FIG. 7 shows a typical XBIOS feature routine organization.

Specific XBIOS Features in the Presently Preferred Embodiment

XBIOS "features" are extensions to standard BIOS that support hardware extensions. A feature is either a hardware routine or data that is embedded into XBIOS. The following list describes various XBIOS features:

- Identify: Identifies the current system
- Setup Entry: Entry point to the ROM based setup program; optional keystroke trigger via ctrl-alt-enter.
- Toggle Speed: Selects the next speed setting; optional keystroke trigger via ctrl-alt-backslash.
- Speed: Set of routines to handle the system speed
 - Returns the number of system speed settings.
 - Returns the current speed setting.
 - Sets the system speed.
- Reverse Video: Reverses the monitor video attributes; optional Keystroke Trigger: via ctrl-alt-backspace.
- Monitor Toggle: Toggles between video monitors; optional keystroke trigger via ctrl-alt-F11.
- Contrast: Set video contrast; optional keystroke trigger via ctrl-alt-F12.
- Shadow RAM: Enable/Disable Shadow Ram
- EMS: Enable/Disable EMS
- Standby: Enable/Disable Standby
- Gate A20: Used to set up fast gate A20.
- Diagnostics: Entry point to memory diagnostic routines.
- Battery: Returns the Battery Voltage Level
- SmartVu: Controls the SmartVu device The SmartVu™ device is a very small character display in

the computer chassis, which is used, under low-level control, to output status and diagnostic messages.

Password: Set/Alter the system password
 Peripheral: Enables/Disables peripheral devices
 Reset: Controls the Reset Button
 Speaker: Controls the speaker volume
 OS2Init: Machine specific initialization for OS/2.

Of course, the disclosed innovative system architecture can be used to add other such features if desired.

Further Modifications and Variations

It will be recognized by those skilled in the art that the innovative concepts disclosed in the present application can be applied in a wide variety of contexts. Moreover, the preferred implementation can be modified in a tremendous variety of ways. Accordingly, it should be understood that the modifications and variations suggested below and above are merely illustrative. These examples may help to show some of the scope of the inventive concepts, but these examples do not nearly exhaust the full scope of variations in the disclosed novel concepts.

For example, the set of extended features can readily be expanded. One way to use this capability is to provide the user with additional debug functions which can be used to interrupt the application software, as desired, to monitor register values, memory usage, etc.

Another advantageous use of the extended feature routines is for dial-up diagnostics (and/or debug). One example of a hardware configuration which is suitable for such dial-up operation is disclosed in published PCT application WO 90/06548, which is hereby incorporated by reference; but of course other hardware configurations can be used instead.

The contemplated primary advantage of the self-describing system software extension feature routines provided by the present invention is for system-operate functions, such as those listed above; but the capabilities of the disclosed architecture can also be exploited advantageously by device drivers for third-party-peripherals. For example, a power-hungry peripheral in a small portable computer can use an XBIOS call to check the battery status before initiating a high-current operation.

The self-describing system software extension feature routines can also be highly advantageous in controlling closely-bundled peripherals. For example, one optional add-on available with most computers in the Dell™ line is a disk drive array controller, known as the Dell Drive Array™ (DDA). The present inventors have already begun work on implementing some control functions for the Dell Drive Array™ with the XBIOS of the presently preferred embodiment, and this direct interface is contemplated as one example of an advantageous use of the disclosed concepts.

For another example, many application developers are struggling with the problem of the range of installed hardware capabilities. Business software may run on a wide range of "IBM-compatible" machines. Even if very old or very low-end machines are excluded, a commercial package such as WordPerfect™ or Paradox™ may be expected to run on anything from an 8-MHz 80286 ISA EGA machine with a crowded 40 msec disk to a 50-MHz 80486 EISA TIGA machine with a disk drive array. This range of machines will provide more than an order of magnitude difference in real-world performance, which poses a dilemma for

application software developers: the features which provide product differentiation, and which run well on high-end machines, will completely bog a lesser machine. Some vendors have responded to this problem by preparing scaled-down versions of their current primary products, to permit operation on machines with less power (such as 8088- or 80286-based portable machines). However, this presents more difficulty in product distribution and support. One way to advantageously exploit the disclosed innovations is for such application software (at installation or startup, or on user command) to use the extended feature table to find out the basic system configuration, and modify its own software configuration or installation accordingly.

Referring to FIG. 8, there is illustrated a flow chart depicting a preferred embodiment of the present invention. At step 800, the CPU initiates operation. Next, at step 801, the CPU retrieves the startup software from nonvolatile memory. Thereafter, at step 802, the CPU automatically begins running the startup software on the CPU including self-test and bootstrap software. Next, at step 803, from execution of the startup software, the CPU launches into execution of the operating system software. At step 804, from execution of the operating system software, the CPU launches into execution of the application software. The CPU, under control of the application software, may programmably call on the basic system software (step 805) from the nonvolatile memory to interface to an I/O device according to a format which is substantially independent of the type of hardware being used within the computer system.

The CPU may also proceed to step 806 where, under control of the application software, it programmably calls on device driver software to interface with an I/O device according to a format which is substantially independent of the computer system hardware. Also, the CPU may proceed to step 807 where, under control of the operating system software it programmably calls on machine-specific system feature extension software to provide a low-level interface to electrical operations. The system feature extension software is partly stored in nonvolatile memory and contains a self-describing feature table and a plurality of machine-dependent routines which are to be executed by the CPU. Note that device driver programs are also able to make calls to the machine-dependent routines which are dependent upon data in the self-describing feature table.

Referring next to FIG. 9, there is illustrated a preferred embodiment of the present invention whereby CPU 900 is coupled to program storage unit 901 from where CPU 900 can read and programmably execute application software programs. CPU 900 is also coupled to I/O devices 902, which include at least one input device and at least one output device. CPU 900 is also coupled to nonvolatile memory (ROM) 903 which contains basic system software at addresses which are accessible by application software programs to provide translation for at least some input and output operations. Startup software stored within nonvolatile memory 903 is called up by CPU 900 whenever CPU 900 initially commences operation. Operation system software configured within either or both nonvolatile memory 903 and volatile memory 904, which is also coupled to CPU 900, is executed by CPU 900 after the startup software is launched. The operating system software allows a user to command the CPU to begin execution of application software programs.

System feature extension software is also stored in nonvolatile memory 903. The system features extension software contains a plurality of machine-dependent routines and a self-describing feature table which contains pointers to the machine-dependent routines. Multiple device driver programs each accessible by the application software programs running on the CPU define a software interface to specific features of at least one of the I/O devices. The device driver programs are able to make calls to the machine-dependent routines which are dependent upon data in the self-describing feature table.

The disclosed innovations have been described with primary reference to a uniprocessor CPU, but they can also be advantageously applied to multiprocessor system.

As will be recognized by those skilled in the art, the innovative concepts described in the present application can be modified and varied over a tremendous range of applications, and accordingly the scope of patented subject matter is not limited by any of the specific exemplary teachings given.

```

;/*****/ ;A03
;/* Check for existence of XBIOS */ ;A03
;/*****/ ;A03
getmaker: ;A03
    mov     cx, XFT_SIGLEN          ;load signature length ;A03
    lea     di, xbhdr.xft_signature ;load addr of signature ;A03
    assume es:dgroup ;A03
    mov     ax, ds ; ;A03
    mov     es, ax ; es -> dgroup ;A03
    assume ds:nothing ;A03
    lds     si, PXFTADDR           ;ds:si -> @ of ptr to xft ;A03
    lds     si, ds:[si]           ;ds:si -> @ of xft ;A03
    repz   cmpsb                  ;check for signature match ;A03
    jz     init_xbios             ;Jump if supported ;A03
    jmp    init_no_xbios          ; jump to old mach_id stuff ;A03
;
; XBIOS exists - use it for initialization
;
init_xbios: ;A03
    mov     xbstatus, 1           ; XBIOS exists ;A03
    add     si, WORDLEN          ; adjust pointer to entries ;A03
    push   si                    ; save as input to find_xbios_ ;A03
    push   ds                    ; ;A03
;
; Locate System Identify Routine
;
    mov     ax, XB_SYSTEM_IDENTIFY ; find system identify entry ;A03
    call    find_xbios_entry ; ;A03
    mov     bx, offset si_data     ; es:bx -> return data area ;A03
    call    ds:[si].xft_ptr       ; call xbios routine ;A03
;
; fill in mach_packet from returned model string and bios revision
;
    push   es ; ;A03
    pop    ds ; ;A03
    mov    di, offset dgroup:mach_packet ; es:di @ mach packet ;A03
    mov    si, offset si_data.system ; ds:si @ of returned string ;A03
xbgetchars: ;A03
    lodsb ; ;A03
    cmp    al, 00h ; copy until 0 encountered ;A03
    jz     xbgetrev ; done, go get bios rev ;A03
    stosb ; ;A03
    loop  xbgetchars ; loop until done ;A03

```

```

xbgetrev:
    mov     di, offset dgroup:mach_packet.rom    ; es:di @ mach packet ;A03
    mov     byte ptr es:[di], "P"              ; Phoenix BIOS ;A03
    inc     di                                  ; point to next char ;A03
xbgetrevchar:
                                                ;A03
    lodsb                                     ; ds:si @ of returned BIOS rev ;A03
    cmp     al, 00h                            ; copy until 0 encountered ;A03
    jz     xbsys_init                          ; done ;A03
    stosb                                     ;A03
    loop   xbgetrevchar                       ; loop until done ;A03
;
; Locate OS/2 init routine
;
xbsys_init:
                                                ;A03
    mov     ax, XB_OS2INIT                     ; find OS2 machine init routine;A03
    call   find_xbios_entry                   ; ;A03
    cmp     ax, 0FFFFh                         ; found? ;A03
    je     smartvu                            ; no - skip machine init ;A03
    call   ds:[si].xft_ptr                    ; call OS2 init routine ;A03
;
; Locate SmartVu routine
;
smartvu:
                                                ;A03
    mov     ax, XB_SMARTVU_ON_OFF             ; find SmartVu On/Off routine ;A03
    call   find_xbios_entry                   ; ;A03
    cmp     ax, 0FFFFh                         ; found? ;A03
    je     CheckPanther                      ; no - skip smartvu init ;A03
;
; Turn On SmartVu (prevent BIOS from overwriting what we do...)
;
    push   es                                 ; save our data segment ;A03
    mov     ax, XB_SMARTVU_ON                 ; turn on SmartVu ;A03
    push   ax                                 ; push as input to XBIOS ;A03
    push   ss                                 ; ;A03
    pop    es                                 ; set es:bx to @ of input parm ;A03
    mov     bx, sp                            ; ;A03
    call   ds:[si].xft_ptr                   ; call SmartVu on/off routine ;A03
    pop    ax                                 ; realign stack ;A03
    pop    es                                 ; restore es ;A03
;
; Write "OS/2" to SmartVu
;
    mov     ax, XB_SMARTVU                   ; find SmartVu Write routine ;A03
    call   find_xbios_entry                   ; ;A03
    cmp     ax, 0FFFFh                         ; found? ;A03
    je     CheckPanther                      ; no - shouldn't happen ;A03
    mov     bx, offset smvu_data              ; es:bx -> data to write ;A03
    call   ds:[si].xft_ptr                   ; call xbios routine ;A03
    jmp    GateA20                            ; go setup Gate A20 ;A03
; Locate Gate A20 XBIOS routine
;

```

```

GateA20:
    mov     ax,XB_A20                ; find system Gate A20 entry ;A03
    call    find_xbios_entry        ;                               ;A03
;
;   set OS/2 gate A20 address to XBIOS A20 routine shell
;
    mov     ax, offset DHA20_xbios   ; offset of gate A20 routine ;A03
    mov     _A20,ax                 ; store it ;A03
;
;   save address of XBIOS GATEA20 routine for real & prot mode
;   (prot mode selector generated on first A20 call - see DHMODESW.ASM)
;
    mov     di, offset fpXbiosA20real ; es:di -> real mode call @ ;A03
    add     si,4                    ; ds:si -> XFT routine address ;A03
    movsw                                     ; store offset ;A03
    movsw                                     ; store segment ;A03
    mov     di, offset fpXbiosA20prot.off ; es:di -> prot mode offset ;A03
    mov     ax, es                   ; restore ds ;A03
    mov     ds, ax                   ;                               ;A03
    mov     si, offset fpXbiosA20real.off ; ds:si -> real mode offset ;A03
    movsw                                     ; store prot mode offset ;A03
;
; Time to leave....
;
    pop     ax                       ; make stack right... ;A03
    pop     ax                       ; (was saved ds:si for find_xbios) ;A03
    jmp     exit_mach_id             ; go to exit ;A03

;*****/;A03
;/* */;A03
;/* ROUTINE NAME: find_xbios_entry */;A03
;/* */;A03
;/* INPUT: */;A03
;/* starting address of XBIOS feature table on stack... */;A03
;/* ax <- XBIOS id to locate */;A03
;/* */;A03
;/* OUTPUT: */;A03
;/* ds:si <- @ of XBIOS feature table entry for id */;A03
;/* ax = -1 (FFFF) means id not found */;A03
;*****/;A03
    assume ds:nothing, es:dgroup ;A03
find_xbios_entry proc near ;A03
    pop     cx                       ; pop return address ;A03
    pop     ds                       ; ds:si -> start of XFT entries ;A03
    pop     si                       ;                               ;A03
    push    si                       ; save on stack for next call... ;A03
    push    ds                       ;                               ;A03
    push    cx                       ; put return address back on stack ;A03
init_feature_loop: ;A03
    cmp     ds:[si].xft_id, XB_NULL   ; are we at the end of XFT? ;A03
    je     init_loop_error           ; yes, break; ERROR ;A03
    cmp     ds:[si].xft_id, XB_CHAIN ; is XFT chain id located? ;A03
    jne    init_found_id             ; if not - we found an id ;A03

```

```

lds    si, ds:[si].xft_ptr      ;set next xft block      ;A03
jmp    SHORT init_feature_loop  ;continue search      ;A03
;                                     ;A03
;* we've found an ID, is it the one we wanted? ;A03
;                                     ;A03
init_found_id: ;A03
    cmp    ds:[si].xft_id,ax      ;is the id the one we want? ;A03
    je     init_loop_exit        ;yes, we're done        ;A03
    mov    bx, XFT_FEATURELEN     ;sizeof(feature) to bx  ;A03
    test   ds:[si].xft_attr, MASK keyb ;is keystroke trigger? ;A03
    jz     init_check_appx      ;if not check for appendix ;A03
    add    bx, XFT_KEYBLEN       ;add length of keyb to bx ;A03
init_check_appx: ;A03
    test   ds:[si].xft_attr, MASK appx ;is there an appendix? ;A03
    jz     init_calc_next_id     ;if not calc next id addr ;A03
    add    bx, ds:[si].[bx]      ;add appendix length to bx ;A03
    add    bx, XFT_APPXLENFLD    ;add size(appx_len_field) ;A03
init_calc_next_id: ;A03
    add    si, bx                ;locate next id        ;A03
    jmp    init_feature_loop     ;continue table scan   ;A03
init_loop_exit: ;A03
    ret                          ;A03
init_loop_error: ;A03
    mov    ax,OFFFh             ;A03
    jmp    init_loop_exit       ;A03
find_xbios_entry endp          ;A03

```

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TITLE PC DOS KEYB Command - Dell feature support

```

;*****
;*                                     *
;* DELLKEYB -- Dell hardware feature support *
;*                                     *
;*                                     *
;* Copyright (C) 1991 Dell Computer Corporation *
;*                                     *
;*****

```

```

    INCLUDE postequ.inc
    EXTRN SHIP_IT:NEAR
    EXTRN mach_id:NEAR
CODE    SEGMENT PUBLIC 'CODE'
        ASSUME CS:CODE,DS:NOTHING
        PUBLIC init_Dell      ;Initialize Dell-specific functions
        PUBLIC int9_swap     ;Temporarily exchange INT9 handlers
        PUBLIC Dell_key      ;Handle ^-alt-keys on all Dell systems
xb_signature    db 'DELLXBIOS',00 ;XBIOS signature string & len
xb_sig_len      equ 10
max_keys        equ 32      ;table sized for n key codes
key_count       dw ?        ;# of codes logged in table
key_codes       dw max_keys dup (?) ;ctl-alt codes handled by XBIOS
key_handlers    dw max_keys*2 dup (?) ;far addresses of XBIOS handlers
current_funct   db ?        ;re-entry avoidance flag
ofs_int9        dw ?        ;old INT9 handler address

```

```

seg_int9      dw ?
setup_F       dd 0f0000100h      ;F000-based ROM setup
setup_E8      dd 0e8000100h      ;E800-based ROM setup
speed_toggle  dd 0f000ff60h      ;ROM-based speed toggle
monitor_toggle dd 0f0000104h      ;ROM-based laptop monitor toggle
contrast_toggle dd 0f000ff86h      ;ROM-based laptop contrast toggle
video_toggle  dd 0f000ff82h      ;ROM-based laptop reverse video toggle

PUBLIC Mach_Type
PUBLIC BIOS_VER
PUBLIC Setup_Addr
PUBLIC LT_System

Mach_Type     db ?
BIOS_VER      db ?
Setup_Addr    db ?
LT_System     db ?

DELL          equ 1              ;machine type
NON_DELL      equ 0
PNX_BIOS      equ 1              ;bios version
PCL_BIOS      equ 0
F_BLOCK       equ 1              ;setup addr
E8_BLOCK      equ 0
DELL_LAPTOP   equ 1              ;laptop flag (LCD monitor)
NOT_DELL_LAPTOP equ 0
ENT_KEY       equ 01CH           ;Enter key used for Setup
BS_KEY        equ 02BH           ;Backslash key for Speed Switch
B_KEY         equ 00Eh           ;backspace key for reverse video.
F11_KEY       equ 057h           ;F11 key for monitor toggle
F12_KEY       equ 058h           ;F12 key for contrast.

CMOS_MC       equ 42h            ;cmos monitor/contrast byte address
MT_FLAG       equ 00001000b      ;monitor toggle flag bit
CE_FLAG       equ 00010000b      ;contrast enhancer flag bit
cmos_data_port equ 71h           ;used for REAL_TIME_CLOCK
cmos_func_port equ 70h           ;used for REAL_TIME_CLOCK
ioj_cmos      equ 8              ;Cmos Ram
iojump        macro dev
               rept   ioj_&dev
               jmp    $+2
               endm
               endm

;-----
in_cmos_data  PROC NEAR          ;AL -> CMOS byte at AL
               pushf
               cli
               or al,80h
               iojump cmos
               out cmos_func_port,al
               iojump cmos
               in al,cmos_data_port
               push ax
               mov al,00h         ;point to register D
               iojump cmos

```

```

out cmos_func_port,al      ;enable NMI
iojump cmos
in al,cmos_data_port      ;requisite read from data port
pop ax
popf
ret
in_cmos_data  ENDP
;*****
swallow_key   PROC NEAR      ;re-enable keyboard
pushf
cli
mov al,20h      ;send end-of-interrupt
out 20h,al
mov al,ENA_KBD  ;re-enable keyboard interrupts for
call SHIP_IT    ;SETUP (DOS 5.0 SHIP_IT doesn't STII)
popf
ret
swallow_key   ENDP
;*****

```

Build Quick Reference Table

The following routine scans the feature table for entities with keystroke activation, and saves a list of the keystroke codes, with the corresponding routine addresses, in system RAM.

This routine can be plugged into a third-party keyboard-handling routine, to preserve the XBIOS keystroke handling when INT9 is remapped, without knowing in advance what the active XBIOS keystroke triggers should be.

```

init_Dell    PROC NEAR

cld
push ax
push bx
push cx
push dx
push bp
push si
push di
push ds
push es

mov cs:key_count,0      ;flag no valid XBIOS keys
mov cs:current_funct,0  ;flag no handlers executing
mov ax,0f000h
mov es,ax
les di,es:[0ed00h]
push ds
push cs

```

```

pop ds
lea si,cs:xb_signature
mov cx,xb_sig_len           ;look for XBIOS signature text
repe cmpsb                 ;at *(F000:ED00)
pop ds
add di,2                   ;skip version ID, if present
jcxz for_feature
go_init_done: jmp xb_init_done           ;no signature, no XBIOS
xbios_chain:  les di,es:[di+2]           ;handle XFT chain feature...
for_feature:  mov ax,es:[di]             ;get feature identifier
add di,2
cmp ax,0                ;NULL feature?
je go_init_done         ;yes, end of feature table
cmp ax,65535            ;CHAIN feature?
je xbios_chain         ;yes, go find next fragment...
mov ax,es:[di]         ;else get feature attributes
add di,2
test ax,00000001b     ;real mode calls accepted?
jz next_feature       ;no, can't call it from KEYB!
test ax,00001000b     ;standard seg:off pointer?
jz next_feature       ;no, wouldn't be prudent at
test ax,00010000b     ;this juncture
jz next_feature       ;not keystroke, don't record
mov bx,es:[di+4]      ;get keystroke data
and bh,11b           ;both ctrl and alt flags set?
xor bh,11b           ;no, key code not accessible
jnz next_feature     ;from standard KEYB.COM
mov si,cs:key_count  ;else record the keystroke in
shl si,1             ;"quick-reference" table...
mov cs:key_codes[si],bx
shl si,1
mov bx,es:[di]
mov cs:key_handlers[si],bx ;write far proc offset
mov bx,es:[di+2]
mov cs:key_handlers[si+2],bx ;write far proc segment
inc cs:key_count      ;bump key counter
cmp cs:key_count,max_keys ;any more room in table?
jae xb_init_done     ;no, stop scanning
next_feature: add di,4 ;skip handler address
test al,110000b     ;keystroke or appendix?
jz for_feature       ;no, nothing else to skip
mov bx,ax
and bx,10000b       ;keystroke?
shr bx,1
shr bx,1            ;BX=2 if keystroke present,
shr bx,1            ;else 0
add di,bx
test ax,100000b     ;appendix?
jz for_feature
add di,es:[di]     ;yes, add blocklen+2 to skip
add di,2
jmp for_feature

```



```

xb_init_done:  mov ax,3509h          ;record current INT9 address
               int 21h           ;so int9_swap can swap it out
               mov seg_int9,es
               mov ofs_int9,bx
               call mach_id       ;identify other machine
                                   ;characteristics

               pop es
               pop ds
               pop di
               pop si
               pop bp
               pop dx
               pop cx
               pop bx
               pop ax
               ret
init_Dell     ENDP

```

```

;*****
int9_swap     PROC NEAR

               push ax
               push bx
               push cx
               push ds
               push es
               xor ax,ax
               mov ds,ax
               les bx,ds:[9*4]    ;pick up current int 9 vector...
               mov ax,cs:ofs_int9 ;and exchange with vector
               mov cx,cs:seg_int9 ;stored during initialization
               mov cs:seg_int9,es
               mov cs:ofs_int9,bx
               pushf
               cli
               mov ds:[9*4],ax
               mov ds:[9*4+2],cx
               popf
               pop es
               pop ds
               pop cx
               pop bx
               pop ax
               ret
int9_swap     ENDP

```

This code uses the Quick Reference Table in system RAM, and provides quasi-real-time handling of keystroke calls.

```

;*****
XBIOS_key    PROC NEAR                ;Check for XBIOS ctrl-alt code in AX
                                           ;Call XBIOS & return C=0 if handled

    cmp cs:current_funcnt,0
    clc                                           ;don't allow re-entry into XBIOS
    jne XB_return                               ;handlers

    mov cs:current_funcnt,1
    push ax
    push bx
    push cx
    push dx
    push bp
    push si
    push di
    push ds
    push es

    mov cx,cs:key_count
    cmp cx,1
    jb XB_exit                                  ;no XBIOS keys logged, exit w/C=1
    xor ah,ah
    push cs
    pop es
    mov di,OFFSET cs:key_codes
    mov dx,di
    cld
    repne scasw                                ;scan key code table
    stc
    jne XB_exit                                  ;key not found, exit w/C=1
    sub di,2                                    ;else point to match
    sub di,dx                                    ;get table offset *2
    shl di,1                                    ;*4 to index jump table
    push di
    call swallow_key
    pop di
    call int9_swap                              ;avoid KEYB re-entrancy
    mov ax,40h                                  ;set DS=40H in case XBIOS wants it
    mov ds,ax
    call DWORD PTR cs:key_handlers[di]
    call int9_swap                              ;put KEYB INT9 handler back
    clc                                           ;signal XBIOS call accomplished

XB_exit:    pop es
            pop ds
            pop di
            pop si
            pop bp
            pop dx

```

```

pop cx
pop bx
pop ax
mov cs:current_funct,0
XB_return:   ret
XBIOS_key   ENDP

```

The following code attempts to call the routine XBIOS_key, to handle XBIOS calls. However, this code can also operate as a simple old-style keystroke handler, for backward compatibility. Comparison of this code with the XBIOS_key routine shows some of the advantages derived from the claimed innovations.

```

;*****
Dell_key    PROC                ;Try to handle key AL through XBIOS
                                           ;(or BIOS on older Dell/PCL machines)

    push ax
    push bx
    push cx
    push dx
    push bp
    push si
    push di
    push ds
    push es
    call XBIOS_key                ;XBIOS function present?
    jnc key_processed            ;yes, keypress handled
    cmp Mach_Type,DELL
    jne not_Dell_key
    cmp BIOS_VER,PNX_BIOS
    jne not_Dell_key            ;not Dell/Phoenix machine

    cmp al,ENT_KEY                ;ENTER triggers ROM-based SETUP
    je call_setup
    cmp al,BS_KEY                ;'\ ' key toggles CPU speed
    je toggle_speed
    cmp LT_System,DELL_LAPTOP
    jne not_Dell_key            ;not a 316LT/320LT, skip other tests
    cmp al,B_KEY                ;backspace key inverts video
    je invert_video
    cmp al,F11_KEY                ;F11 toggles LCD/external monitor
    je toggle_monitor
    cmp al,F12_KEY                ;F12 toggles enhanced LCD contrast
    je toggle_contrast
    jmp short not_Dell_key

key_processed: db 85h                ;stc becomes test cx,di; sets C=0
not_Dell_key:  stc                ;key not handled in BIOS

    pop es
    pop ds
    pop di
    pop si

```

```

    pop bp
    pop dx
    pop cx
    pop bx
    pop ax
    ret

call_setup:                                ;Invoke ROM-based SETUP program
    call swallow_key
    call int9_swap                          ;ditch the KEYB keypress handler
    cmp Setup_Addr,F_BLOCK
    jne E8_setup
    call [setup_F]                          ;far call to SETUP vector
    jmp short end_setup
E8_setup:    call [setup_E8]
end_setup:  call int9_swap                    ;restore our INT9 handler
    jmp key_processed

toggle_speed:                                ;Toggle CPU speed
    call [speed_toggle]
    jmp key_processed

invert_video:                                ;Invert LCD video
    call [video_toggle]
    jmp key_processed

toggle_monitor:                              ;Select LCD/external monitor
    mov al,CMOS_MC
    call in_cmos_data                       ;...but only if option is active
    test al,MT_FLAG
    jz not_Dell_key                         ;else ignore key
    call [monitor_toggle]
    jmp key_processed

toggle_contrast:                             ;Toggle LCD contrast enhancement
    mov al,CMOS_MC
    call in_cmos_data                       ;...but only if option is active
    test al,CE_FLAG
    jz not_Dell_key                         ;else ignore key
    call [contrast_toggle]
    jmp key_processed

Dell_key    ENDP
;*****
CODE        ENDS
           END

```



```

les bx,es:[di+2]
mov WORD PTR lpExtA20Handler,bx ;else record handler addr...
mov WORD PTR lpExtA20Handler+2,es
stc                               ;...and return OK
xb_init_done: pop es                ;C=1 if XBIOS OK, 0 otherwise
pop ds
popa
mov ax,0
rcl ax,1
popf
ret
next_feature: mov ax,es:[di]        ;get feature attributes
add di,6                          ;skip attribs & handler addr
test al,110000b                   ;keystroke or appendix?
jz for_feature                    ;no, nothing else to skip
mov bx,ax
and bx,10000b                     ;keystroke?
shr bx,1
shr bx,1                          ;BX=2 if keystroke present,
shr bx,1                          ;else 0
add di,bx
test ax,100000b                   ;appendix?
jz for_feature
add di,es:[di]                   ;yes, add blocklen+2 to skip
add di,2
jmp for_feature
IsDellXBIOS   ENDP

```

```

;* XBIOS.INC *;
;* ----- *;
; $Log: X:/bios/core/xbios/xbios.inv $
;* ----- *;
comment ;*
    Include file contains the XBIOS structure declarations for the header
    and the permanent part of the feature entry. Also included are bit
    field records for flag entries and equates for the feature entry id's.
;*
;* XFT Header Structure. *;
xft_header_tag    struc
    xft_signature  db "DELLXBIOS",0
    xft_version    dw 0100h
xft_header_tag    ends

;* XFT Feature Structure. *;
xft_feature_tag   struc
    xft_id         dw ?
    xft_attr       dw ?
    xft_ptr        dd ?
xft_feature_tag   ends

;* XFT Keystroke Trigger Structure. *;
xft_keyb_tag      struc
    xft_keyb       dw 0

```

```

xft_keyb_tag      ends
;* XFT Appendix Structure.
xft_appx_tag      struc
    xft_appx       dw 0
    xft_appx_dat   db ?
xft_appx_tag      ends
;
;* XFT POINTER LOCATION (Offset into F blk).
XFT_PTR_LOC       equ    0ED00h
XFT_LINR_PTR_LOC  equ    0FED00h
;* XFT LENGTH EQUATES
XFT_HDRLEN        equ    (size xft_header_tag)
XFT_FEATURELEN    equ    (size xft_feature_tag)
XFT_SIGNLEN       equ    (offset xft_version - offset xft_signature)
XFT_KEYBLEN       equ    (size xft_keyb_tag)
XFT_APPXLENFLD    equ    (size xft_appx)
;* XFT FEATURE ATTRIBUTE FLAGS.
XATTR_REAL        equ    0000000000000001b ;xftptr to real mode proc
XATTR_PROT        equ    0000000000000010b ;xftptr to protect mode proc
XATTR_XXXX        equ    0000000000000100b ;reserved bit
XATTR_SEGM        equ    0000000000001000b ;xftptr in seg:off format
XATTR_KEYB        equ    0000000000010000b ;optional keystroke trigger
XATTR_APPX        equ    0000000000100000b ;optional appendix
XATTR_BIMODAL     equ    (XATTR_REAL or XATTR_PROT)
XATTR_DATA        equ    (XATTR_BIMODAL xor XATTR_BIMODAL)
XATTR_LINR        equ    (XATTR_SEGM xor XATTR_SEGM)
XFT_ATTR_FLAGS    record  resv:10=0,appx:1,keyb:1,form:1,xxxx:1=0,prot:1,real:1
;* XFT KEYBOARD TRIGGER SHIFT STATE FLAGS.
XKEYB_CTRL        equ    00000001b          ;ctrl key
XKEYB_ALT         equ    00000010b         ;alt key
XFT_KEYB_FLAGS    record  unused:6=0,alt:1,ctrl:1
;
;* -----
;* XBIOS FEATURE EQUATES
comment ;*
    The following are equates used to identify XBIOS extended features.
    The XB_NULL and XB_CHAIN entries are reserved. Place new entries before
    the XB_CHAIN and include a short description of the feature and its
    protocol.
;*
;* -----
comment ;*
    * The NULL feature entry indicates the end of the extended feature table.
;*
XB_NULL           equ    00000h ;XFT termination identifier
;* -----
comment ;*
    * FUNCTION: System Identify
    * INPUT: es:bx - points to output stack frame
    *         cs,ds - XBIOS routine segment/selector (must be F000 block)
    * INPUT STACK FRAME: None

```

```

* OUTPUT: Results are put in output stack frame, which is treated
*         like a buffer. For long names, could be up to 32 bytes long.
*         The first byte stored in the buffer is the Dell system
*         model number (binary), followed by the Dell system revision
*         number (binary). Next is the system name, which is an
*         ASCII string terminated with a zero. Finally, the BIOS
*         version is included, another ASCII string terminated with
*         a zero.
* OUTPUT STACK FRAME:
*     system_id  db 32 dup(?)
;*
XB_SYSTEM_IDENTIFY  equ          00001h
;* ----- *;
comment ;*
* FUNCTION:
* INPUT:
* INPUT STACK FRAME:
* OUTPUT:
* OUTPUT STACK FRAME:
;*
XB_SETUP_ENTRY      equ          00002h
;* ----- *;
comment ;*
* FUNCTION: Toggle speed
* INPUT: None
* INPUT STACK FRAME: No input
* OUTPUT: None
* OUTPUT STACK FRAME: No output
;*
XB_TOGGLE_SPEED     equ          00003h
;* ----- *;
comment ;*
* FUNCTION: Speed Control and Status
* INPUT: es:bx - points to input stack frame
*         cs   - XBIOS routine segment/selector (must be F000 block)
* INPUT STACK FRAME:
* OUTPUT:
* OUTPUT STACK FRAME:
;*
XB_SPEED_CONTROL    equ          00004h
;* ----- *;
comment ;*
* FUNCTION: Toggle black/white video background on portable systems
* INPUT: None
* INPUT STACK FRAME: No input
* OUTPUT: None
* OUTPUT STACK FRAME: No output
;*
XB_REVERSE_VIDEO    equ          00005h
;* ----- *;
comment ;*
* FUNCTION: Toggle CRT/LCD displays on portable systems

```



```

* INPUT: None
* INPUT STACK FRAME: No input
* OUTPUT: None
* OUTPUT STACK FRAME: No output
;*
XB_MONITOR_TOGGLE equ 00006h
;* -----*
comment ;*
* FUNCTION: Toggle text mode contrast / graphics mode palettes on portables
* INPUT: None
* INPUT STACK FRAME: None
* OUTPUT: None
* OUTPUT STACK FRAME: None
;*
XB_CONTRAST equ 00007h
;* -----*
comment ;*
* FUNCTION:
* INPUT:
* INPUT STACK FRAME:
* OUTPUT:
* OUTPUT STACK FRAME:
;*
XB_SHADOW_RAM equ 00008h
;* -----*
comment ;*
* FUNCTION:
* INPUT:
* INPUT STACK FRAME:
* OUTPUT:
* OUTPUT STACK FRAME:
;*
XB_EMS equ 00009h
;* -----*
comment ;*
* FUNCTION: Disable/enable "Standby" keyboard NMI on portable systems
* INPUT: es:bx - points to input stack frame
* INPUT STACK FRAME: First byte 0 to disable standby key, 1 to enable
* OUTPUT: None
* OUTPUT STACK FRAME: None
;*
XB_STANDBY equ 0000Ah
XB_STANDBY_DISABLE equ 0
XB_STANDBY_ENABLE equ 1
;* -----*
comment ;*
* FUNCTION: Enable/Disable A20 Line
* INPUT: es:bx - points to input stack frame
* cs - XBIOS routine segment/selector (must be F000 block)
* The input stack frame contains the subfunction id:
* - 0 means disable A20 line (wrap addresses at 1MB)
* - 1 means enable A20 line (allow high memory access)

```

```

* INPUT STACK FRAME:
*       subfunc db ?
* OUTPUT: A20 line enabled/disabled
* OUTPUT STACK FRAME: None
;*
XB_A20          equ      0000Bh
XB_A20_DISABLE  equ      0
XB_A20_ENABLE   equ      1
;* ----- *;
comment ;*
* FUNCTION:
* INPUT:
* INPUT STACK FRAME:
* OUTPUT:
* OUTPUT STACK FRAME:
;*
XB_DIAGNOSTICS  equ      0000Ch
;* ----- *;
comment ;*
* FUNCTION:
* INPUT:
* INPUT STACK FRAME:
* OUTPUT:
* OUTPUT STACK FRAME:
;*
XB_BATTERY      equ      00000h
;* ----- *;
comment ;*
* FUNCTION: Output characters on SmartVu device
* INPUT: es:bx - points to input stack frame
*       cs  - XBIOS routine segment/selector (must be F000 block)
*       The stack frame contains the ASCIIZ string to be displayed
*       on the SmartVu device. es:bx points to the first character
*       (leftmost) to be displayed.
* INPUT STACK FRAME:
*       string db x dup (?) - ASCIIZ string of any length
* OUTPUT: ASCIIZ string displayed on SmartVu device
* OUTPUT STACK FRAME: None
;*
XB_SMARTVU      equ      0000Eh
;* ----- *;
comment ;*
* FUNCTION:
* INPUT:
* INPUT STACK FRAME:
* OUTPUT:
* OUTPUT STACK FRAME:
;*
XB_PASSWORD     equ      0000Fh
;* ----- *;
comment ;*
* FUNCTION:

```

```

* INPUT:
* INPUT STACK FRAME:
* OUTPUT:
* OUTPUT STACK FRAME:
;*
XB_PERIPHERAL      equ      00010h
;* ----- *;
comment ;*
* FUNCTION:
* INPUT:
* INPUT STACK FRAME:
* OUTPUT:
* OUTPUT STACK FRAME:
;*
XB_RESET           equ      00011h
;* ----- *;
comment ;*
* FUNCTION: Perform OS/2 Specific Machine Initialization
* INPUT:  cs - XBIOS routine segment/selector (must be F000 block)
* INPUT STACK FRAME: None
* OUTPUT: machine specific actions performed, no data returned
* OUTPUT STACK FRAME: None
;*
XB_OS2INIT        equ      00012h
;* ----- *;
comment ;*
* FUNCTION: Turn SmartVu On/Off
* INPUT:  es:bx - points to input stack frame
*         cs   - XBIOS routine segment/selector (must be F000 block)
*         The input stack frame contains the subfunction id:
*         - 0 turns on SmartVu (disables BIOS writes to SmartVu)
*         - 1 turns off SmartVu (enables BIOS writes to SmartVu)
* INPUT STACK FRAME:
*         subfunc db ?
* OUTPUT: BIOS updates to SmartVu enabled or disabled
* OUTPUT STACK FRAME: None
;*
XB_SMARTVU_ON_OFF equ      00013h
XB_SMARTVU_ON     equ      0
XB_SMARTVU_OFF    equ      1
;* ----- *;
comment ;*
* FUNCTION: Turn system processor cache on/off
* INPUT:  es:bx - points to input stack frame
*         cs   - XBIOS routine segment/selector (must be F000 block)
* INPUT STACK FRAME:
*         status db ?
* OUTPUT: Machine specific actions performed, no data returned
* OUTPUT STACK FRAME:
;*

```

```

XB_SYSTEM_CACHE      equ      00014h
XB_SYSCACHE_OFF      equ      0
XB_SYSCACHE_ON       equ      1
;* -----*
comment ;*
* FUNCTION:   Unlock the FE3021 chip
* INPUT:      cs, ds - selector to 0F000h
* INPUT STACK FRAME:
* OUTPUT:     ZF = 1 : Unlock failure
*             ZF = 0 : Unlock O.K
* OUTPUT STACK FRAME:
;*
xb_unlock            equ      00015h
;* -----*
comment ;*
* FUNCTION:   Get/Set Monitor Type, System Video Status
* INPUT:      es:bx - points to input stack frame
*             cs   - XBIOS routine segment/selector (must be F000 block)
* INPUT STACK FRAME:
* OUTPUT:
* OUTPUT STACK FRAME:
;*
XB_MONITOR_TYPE      equ      00016h
;* -----*
comment ;*
    The chain feature entry indicates that the XFT is fragmented and that
    its corresponding XFT pointer references the next fragment.
;*
XB_CHAIN              equ      0FFFFh ;XFT chain indicator.
;* -----*
;* -----*
;* -----*

file XBIOS . H      sample XBIOS table definitions in C
//
typedef struct
{
    unsigned model_num;
    unsigned board_rev;
    char name[32];
    char bios_version[16];
} sysID;
unsigned get_XFT_version(void);
char far *find_feature(unsigned feature_id);
unsigned system_identify(sysID *ID);
unsigned get_system_speed(unsigned *speed);
unsigned get_supported_speed_count(unsigned *speed);
unsigned set_system_speed(unsigned speed);
unsigned set_A20(unsigned status);
unsigned set_system_cache(unsigned status);
unsigned rom_setup(void);
unsigned reverse_video(void);
unsigned monitor_toggle(void);

```

```

unsigned speed_toggle(void);
unsigned contrast_toggle(void);
unsigned set_standby(unsigned status);
unsigned OS2_init(void);
unsigned write_smartvu(char far *string);
unsigned set_smartvu(unsigned status);
unsigned unlock_3021(void);
unsigned get_monitor_type(unsigned *montype);
unsigned set_monitor_type(unsigned montype);
unsigned get_system_video_status(unsigned *status);
#define XB_SIG_LEN 10
#define XB_NULL 0
#define XB_CHAIN 65535
#define XB_SYSTEM_IDENTIFY 1
#define XB_SETUP_ENTRY 2
#define XB_TOGGLE_SPEED 3
#define XB_SPEED_CONTROL 4
#define XB_REVERSE_VIDEO 5
#define XB_MONITOR_TOGGLE 6
#define XB_CONTRAST 7
#define XB_SHADOW_RAM 8
#define XB_EMS 9
#define XB_STANDBY 10
#define XB_A20 11
#define XB_DIAGNOSTICS 12
#define XB_BATTERY 13
#define XB_SMARTVU 14
#define XB_PASSWORD 15
#define XB_PERIPHERAL 16
#define XB_RESET 17
#define XB_OS2INIT 18
#define XB_SMARTVU_ON_OFF 19
#define XB_SYSTEM_CACHE 20
#define XB_UNLOCK_3021 21
#define XB_MONITOR_TYPE 22

```

```

file XAPI    C    Sample XBIOS test code in C
/*****/
/*          */
/* XAPI -- Turbo C++ XBIOS application program interface */
/*          */
/* Compile with Turbo C++ v1.0 or later, medium model */
/*          */
/*          */
/*****/
#pragma inline
#include <stdio.h>
#include <dos.h>
#include <stdlib.h>
#include <alloc.h>
#include "xbios.h"
/*****/

```

```

unsigned get_XFT_version(void)
(
    unsigned XFT_vers;
    static char xb_signature[] = ("DELLXBIOS");
    asm(
        mov ax,0f000h; mov es,ax; les di,es:[0ed00h];
        lea si,xb_signature; mov cx,XB_SIG_LEN; cld; repe cmpsb;
        jne not_valid;
        mov ax,es:[di]; mov XFT_vers,ax;
    )
    return XFT_vers;
not_valid:
    return 0;
)
/*****/
char far *find_feature(unsigned feature_id)
(
    char far *pntr = NULL;
    asm(
        mov ax,0f000h; mov es,ax; les di,es:[0ed00h];
        add di,XB_SIG_LEN+2;
    )
cmp_loop:
    asm(
        mov ax,es:[di]; add di,2;
        cmp ax,feature_id; je found;
        cmp ax,XB_NULL; je eot;
        mov ax,es:[di]; add di,6;
        test al,110000b; jz cmp_loop;
        mov bx,ax; and bx,10000b; shr bx,1; shr bx,1; shr bx,1; add di,bx;
        test al,100000b; jz cmp_loop;
        mov ax,[di]; add di,2; add di,ax; jmp cmp_loop;
    )
found:
    asm(
        mov ax,es:[di]; add di,2;
        and ax,1000b; jz linear_ptr;
        les di,es:[di];
        mov WORD PTR pntr,di; mov WORD PTR pntr+2,es;
        jmp eot;
    )
linear_ptr:
    asm(
        mov ax,es:[di+2]; mov ah,al; xor al,al; mov cl,4; shl ax,cl; mov es,ax;
        mov di,es:[di];
        mov WORD PTR pntr,di; mov WORD PTR pntr+2,es;
        jmp eot;
    )
eot:
    return pntr;
)
/*****/

```

```

unsigned system_identify(sysID *ID)
{
    unsigned char buffer[32];
    char far *feature;
    if ((feature=find_feature(XB_SYSTEM_IDENTIFY)) == NULL)
        return 0;
    asm(
        push ss; pop es; lea bx,buffer;
        push ds; push si; push di; push bp; mov ax,0f000h; mov ds,ax;
        call DWORD PTR [feature];
        pop bp; pop di; pop si; pop ds;
    )
    ID->model_num=buffer[0];
    ID->board_rev=buffer[1];
    strcpy(ID->name,&buffer[2]);
    strcpy(ID->bios_version,&buffer[3+strlen(&buffer[2])]);
    return 1;
}
/*****/
unsigned get_system_speed(unsigned *speed)
{
    char far *feature;
    unsigned char fn[4];
    if ((feature=find_feature(XB_SPEED_CONTROL)) == NULL)
        return 0;
    fn[3] = 0;
    asm(
        push ss; pop es; lea bx,fn;
        push ds; push si; push di; push bp;
        mov ax,40h; mov ds,ax;
        call DWORD PTR [feature];
        pop bp; pop di; pop si; pop ds;
        cmp ax,0; jne error;
    )
    *speed = fn[2];
    return 1;
error:
    return 0;
}
/*****/
unsigned get_supported_speed_count(unsigned *speed)
{
    char far *feature;
    unsigned char fn[4];

    if ((feature=find_feature(XB_SPEED_CONTROL)) == NULL)
        return 0;
    fn[3] = 0;
    asm(
        push ss; pop es; lea bx,fn;
        push ds; push si; push di; push bp;
        mov ax,40h; mov ds,ax;

```

```

    call DWORD PTR [feature];
    pop bp; pop di; pop si; pop ds;
    cmp ax,0; jne error;
    )
    *speed = fn[1];
    return 1;
error:
    return 0;
}
/*****/
unsigned set_system_speed(unsigned speed)
{
    char far *feature;
    unsigned char fn[4];

    if ((feature=find_feature(XB_SPEED_CONTROL)) == NULL)
        return 0;
    fn[3] = 1;
    fn[0] = speed;
    asm(
        push ss; pop es; lea bx,fn;
        push ds; push si; push di; push bp;
        mov ax,40h; mov ds,ax;
        call DWORD PTR [feature];
        pop bp; pop di; pop si; pop ds;
        cmp ax,0; jne error;
    )
    return 1;
error:
    return 0;
}
/*****/
unsigned set_A20(unsigned status)
{
    char far *feature;

    if ((feature=find_feature(XB_A20)) == NULL)
        return 0;
    asm(
        push ss; pop es; lea bx,status;
        push ds; push si; push di; push bp;
        call DWORD PTR [feature];
        pop bp; pop di; pop si; pop ds;
    )
    return 1;
}
/*****/
unsigned set_system_cache(unsigned status)
{
    char far *feature;

    if ((feature=find_feature(XB_SYSTEM_CACHE)) == NULL)
        return 0;

```



```

asm(
    push es; pop es; lea bx,status;
    push ds; push si; push di; push bp;
    call DWORD PTR [feature];
    pop bp; pop di; pop si; pop ds;
)
return 1;
)
/*****/
unsigned rom_setup(void)
(
    char far *feature;

    if ((feature=find_feature(XB_SETUP_ENTRY)) == NULL)
        return 0;
    asm(
        push ds; push si; push di; push bp;
        call DWORD PTR [feature];
        pop bp; pop di; pop si; pop ds;
    )
    return 1;
)
/*****/
unsigned reverse_video(void)
(
    char far *feature;

    if ((feature=find_feature(XB_REVERSE_VIDEO)) == NULL)
        return 0;
    asm(
        push ds; push si; push di; push bp;
        call DWORD PTR [feature];
        pop bp; pop di; pop si; pop ds;
    )
    return 1;
)
/*****/
unsigned monitor_toggle(void)
(
    char far *feature;

    if ((feature=find_feature(XB_MONITOR_TOGGLE)) == NULL)
        return 0;
    asm(
        push ds; push si; push di; push bp;
        call DWORD PTR [feature];
        pop bp; pop di; pop si; pop ds;
    )
    return 1;
)
/*****/
unsigned contrast_toggle(void)
(

```

```

char far *feature;

if ((feature=find_feature(XB_CONTRAST)) == NULL)
    return 0;
asm(
    push ds; push si; push di; push bp;
    call DWORD PTR [feature];
    pop bp; pop di; pop si; pop ds;
)
return 1;
)
/*****/
unsigned speed_toggle(void)
(
    char far *feature;

    if ((feature=find_feature(XB_TOGGLE_SPEED)) == NULL)
        return 0;
    asm(
        push ds; push si; push di; push bp;
        call DWORD PTR [feature];
        pop bp; pop di; pop si; pop ds;
    )
    return 1;
)
/*****/
unsigned set_standby(unsigned status)
(
    char far *feature;

    if ((feature=find_feature(XB_STANDBY)) == NULL)
        return 0;
    asm(
        push es; pop es; lea bx,status;
        push ds; push si; push di; push bp;
        call DWORD PTR [feature];
        pop bp; pop di; pop si; pop ds;
    )
    return 1;
)
/*****/
unsigned OS2_init(void)
(
    char far *feature;

    if ((feature=find_feature(XB_OS2INIT)) == NULL)
        return 0;
    asm(
        push ds; push si; push di; push bp;
        call DWORD PTR [feature];
        pop bp; pop di; pop si; pop ds;
    )
)

```

```

    return 1;
}
/*****/
unsigned set_smartvu(unsigned status)
{
    char far *feature;

    if ((feature=find_feature(XB_SMARTVU_ON_OFF)) == NULL)
        return 0;
    asm(
        push es; pop es; lea bx,status;
        push ds; push si; push di; push bp;
        call DWORD PTR [feature];
        pop bp; pop di; pop si; pop ds;
    )
    return 1;
}
/*****/
unsigned write_smartvu(char far *string)
{
    char far *feature;

    if ((feature=find_feature(XB_SMARTVU)) == NULL)
        return 0;
    asm(
        les bx,string;
        push ds; push si; push di; push bp;
        call DWORD PTR [feature];
        pop bp; pop di; pop si; pop ds;
    )
    return 1;
}
/*****/
unsigned unlock_3021(void)
{
    char far *feature;

    if ((feature=find_feature(XB_UNLOCK_3021)) == NULL)
        return 0;
    asm(
        push ds; push si; push di; push bp;
        mov ax,0f000h; mov ds,ax;
        call DWORD PTR [feature];
        pop bp; pop di; pop si; pop ds;
    )
    return 1;
}
/*****/
unsigned get_monitor_type(unsigned *montype)
{
    char far *feature;
    unsigned char fn[4];

```

```

if ((feature=find_feature(XB_MONITOR_TYPE)) == NULL)
    return 0;
fn[3] = 0;
asm(
    push ss; pop es; lea bx,fn;
    push ds; push si; push di; push bp;
    mov ax,40h; mov ds,ax;
    call DWORD PTR [feature];
    pop bp; pop di; pop si; pop ds;
    cmp ax,0; jne error;
)
*montype = fn[2];
return 1;
error:
    return 0;
)
/*****/
unsigned get_system_video_status(unsigned *status)
(
    char far *feature;
    unsigned char fn[4];

    if ((feature=find_feature(XB_MONITOR_TYPE)) == NULL)
        return 0;
    fn[3] = 0;
    asm(
        push ss; pop es; lea bx,fn;
        push ds; push si; push di; push bp;
        mov ax,40h; mov ds,ax;
        call DWORD PTR [feature];
        pop bp; pop di; pop si; pop ds;
        cmp ax,0; jne error;
    )
    *status = fn[1];
    return 1;
error:
    return 0;
)
/*****/
unsigned set_monitor_type(unsigned montype)
(
    char far *feature;
    unsigned char fn[4];
    if ((feature=find_feature(XB_MONITOR_TYPE)) == NULL)
        return 0;
    fn[3] = 1;
    fn[0] = montype;
    asm(
        push ss; pop es; lea bx,fn;
        push ds; push si; push di; push bp;
        mov ax,40h; mov ds,ax;
        call DWORD PTR [feature];

```

```

    pop bp; pop di; pop si; pop ds;
    cmp ax,0; jne error;
  }
  return 1;
error:
  return 0;
}

```

What is claimed is:

1. A method for operating a computer system, comprising the steps of:

- (a) automatically running startup software, from non-volatile memory, on a central processing unit (CPU) whenever said computer system first initiates operation, said startup software including self-test and bootstrap software;
- (b) launching said CPU, from execution of said startup software, into execution of operating system software; and
- (c) launching said CPU, from execution of said operating system software, into execution of application software;
- (d) wherein said CPU, under control of said application software, programmably calls on one of
 - (d1) said operating system software, for interfacing to an I/O device according to a format which is substantially independent of a hardware type of said computer system, said CPU, under control of said operating system software, programmably calling machine-specific system feature extension software to provide low-level interface to electrical operations, said system feature extension software being partly stored in said non-volatile memory and containing a self-describing feature table and a plurality of machine-dependent routines executed by said CPU; and
 - (d2) device driver software, stored in said nonvolatile memory, for interfacing to said I/O device according to a format which is substantially independent of said hardware type of said computer system.

2. The method of claim 1, wherein said basic system software is interrupt-driven.

3. The method of claim 1, wherein said device driver programs include a keyboard interface driver program.

4. The method of claim 1, wherein said device driver programs includes a modem interface driver program.

5. The method of claim 1, wherein said device driver programs include a graphics-interface-card driver program.

6. The method of claim 1, wherein said device driver programs includes a mouse interface driver program.

7. The method of claim 1, wherein said self-describing feature table includes plural items, said plural items including callable routines and readable data structures.

8. The method of claim 1, wherein said nonvolatile memory is an EPROM.

9. The method of claim 1, wherein said operating system software runs continuously as a background process on said CPU.

10. The method of claim 1, wherein said basic system software provides translation for input and output operations to a mass storage device.

11. The method of claim 1, wherein said basic system software and said self-describing feature table are both stored in a ROM.

12. The method of claim 1, wherein said basic system software is stored in a ROM, and said ROM also contains a pointer to an address of said self-describing feature table, and said self-describing feature table is stored in rewritable nonvolatile memory.

13. The method of claim 1, wherein said device driver software also provides translation for input and output operations to a peripheral device which shares a data bus with said CPU.

14. The method of claim 1, wherein said step of launching said CPU into execution of operating system software is performed automatically.

15. A method for operating a computer system which includes a central processing unit (CPU), a program storage unit connected to said CPU so that said CPU can read and programmably execute application software programs therefrom, plural I/O devices including an input device and an output device, a mass storage device, and a nonvolatile memory, said method comprising the steps of:

(a) automatically running startup software on said CPU whenever said system first initiates operation, said startup software including self-test and bootstrap software, said startup software being retrieved from nonvolatile memory;

(b) launching said CPU, from execution of said startup software, into execution of operating system software; and

(c) launching said CPU, from execution of said operating system software, into execution of application software;

(d) wherein said CPU, under control of said application software, programmably calls on one of:

- (d1) said operating system software, for interfacing to a selected one of said I/O devices according to a format which is substantially independent of a hardware type of said computer system, said CPU, under control of said operating system software, calling on machine-specific system feature extension software to provide low-level interface to electrical operations, said system feature extension software being partly stored in said nonvolatile memory and containing a self-describing feature table and a plurality of machine-dependent routines executed by said CPU; and
- (d2) device drive software, stored in said nonvolatile memory, for interfacing to said selected I/O device according to a format which is substantially independent of said hardware type of said computer system.

16. The method of claim 15, wherein said basic system software is interrupt-driven.

17. The method of claim 15, wherein said device driver programs include a keyboard interface driver program.

18. The method of claim 15, wherein said self-describing feature table includes plural items, said plural items including callable routines and readable data structures.

19. The method of claim 15, wherein said basic system software provides translation for input and output operations to said mass storage device.

20. The method of claim 15, wherein said basic system software and said self-describing feature table are both stored in a ROM.

21. The method of claim 15, wherein said basic system software is stored in a ROM, and said ROM also contains a pointer to an address of said self-describing feature table, and said self-describing feature table is stored in rewritable nonvolatile memory.

22. The method of claim 15, wherein said device driver software also provides translation for input and output operations to a peripheral device which shares a data bus with said CPU.

23. The method of claim 15, wherein said step of launching said CPU into execution of operating system software is performed automatically.

24. A computer system, comprising:

a central processing unit (CPU);

a program storage unit, coupled to said CPU so that said CPU can programmably execute application software therefrom;

plural I/O devices coupled to said CPU, including an input device and an output device;

operating system software, running on said CPU, and accessible by said application software running on said CPU;

multiple device driver programs, each accessible by said application software running on said CPU to define a software interface to specific features of at least one said I/O device; and

nonvolatile memory coupled to said CPU, containing a machine-specific system feature extension software to provide a low-level interface to electrical operations, said system feature extension software containing a plurality of machine-dependent routines, and a self-describing feature table which contains pointers to said machine-dependent routines,

wherein said device driver programs can make calls to said machine-dependent routines which are dependent on data in said self-describing feature table, and

wherein said CPU, under control of said application software, programmably calls on one of:

said operating system software, for interfacing to a selected one of said plural I/O devices according to a format which is substantially independent of a hardware type of said computer system, said CPU, under control of said operating system software, calling on said machine-dependent system feature extension software to provide said low-level interface to electrical operations; and

said device driver programs for interfacing to said selected I/O device according to a format which is substantially independent of said hardware type of said computer system.

25. The system of claim 24, wherein said application software programs can make calls directly to said self-

describing feature table and thereby to a desired one of said machine-dependent routines.

26. The system of claim 24, wherein said self-describing feature table includes plural items, said plural items including callable routines and readable data structures.

27. The system of claim 24, wherein at least one said input device is a keyboard.

28. The system of claim 24, wherein at least one said output device is a display.

29. The system of claim 24, further comprising installed graphical-user-interface software which adds to the functionality of said operating system software.

30. The system of claim 24, wherein said operating system software is also configured so that said CPU automatically returns to execution of said operating system software after normal termination of any top-level program of the application software.

31. The system of claim 24, wherein said operating system software is stored at an address, in said nonvolatile memory, which is accessible to said CPU.

32. The system of claim 24, wherein said operating system software is stored at an address, on a magnetic recording medium, which is accessible to said CPU.

33. The system of claim 24, wherein said operating system software runs continuously as a background process on said CPU.

34. The system of claim 24, wherein at least one said program storage unit is RAM.

35. The system of claim 24, wherein at least one said program storage unit comprises both a nonvolatile rewritable mass storage medium and a random-access memory.

36. The system of claim 24, wherein at least one said program storage unit consists of a volatile RAM.

37. The system of claim 24, further comprising startup software running on said CPU which includes a system self-test routine and a bootstrap program-loading routine.

38. A computer system, comprising:

a central processing unit (CPU);

at least one program storage unit, coupled to said CPU so that said CPU can programmably execute application software programs therefrom;

plural I/O devices coupled to said CPU, including at least one input device and at least one output device;

a nonvolatile memory coupled to said CPU containing basic system software at addresses which are accessible by said application software programs to provide translation for at least some input and output operations;

startup software, stored in said nonvolatile memory, which is connected so that said CPU call calls said startup software whenever said CPU initially commences operation;

operating system software, configured within computer system memory so that said startup software launches said CPU into execution of said operating system software, and so that a user can command said CPU, through said operating system software, to begin execution of an application software program, and

system feature extension software, stored in said nonvolatile memory, which contains a plurality of machine-dependent routines, and a self-describing

feature table which contains pointers to said machine-dependent routines;
 wherein said CPU, under control of said application software program, programmably calls on one of: said operating system software, for interfacing to a selected one of said plural I/O devices according to a format which is substantially independent of a hardware type of said computer system, said CPU, under control of said operating system software, calling on said system feature extension software to provide a low-level interface to electrical operations; and
 device driver software for interfacing to said selected I/O device according to a format which is substantially independent of said hardware type of said computer system.

39. The system of claim 38, wherein said basis system software provides translation for input and output operations to at least one mass storage device.

40. The system of claim 38, wherein said basis system software is stored in ROM.

41. The system of claim 38, wherein said basis system software is stored in a ROM, and said ROM also contains a pointer to the address of said self-describing feature table.

42. The system of claim 38, wherein said basic system software is stored in a ROM, and said ROM also contains a pointer to the address of said self-describing feature table, and said self-describing feature table is stored in rewritable nonvolatile memory.

43. The system of claim 38, wherein said basic system software and said self-describing feature table are both stored in a ROM.

44. The system of claim 38, wherein said device driver software can make calls to said machine-dependent routines, which are dependent on data in said self-describing feature table.

45. The system of claim 38, wherein said device driver software provides translation for input and output operations to at least one peripheral device which shares a data bus with said CPU.

46. A family of computer systems, said family comprising first, second, and third pluralities of systems; wherein each individual one of said computer systems comprises:
 a central processing unit (CPU);
 a program storage unit, coupled to said CPU so that said CPU can programmably execute application software programs therefrom;
 plural I/O devices coupled to said CPU, including an input device and an output device;
 a nonvolatile memory coupled to said CPU containing basis system software at addresses which are accessible by application software programs to provide translation for at least some input and output operations;

startup software, stored in said nonvolatile memory, which is connected so that said CPU calls said

startup software whenever said CPU initially commences operation;
 operating system software, configured within computer system memory so that said startup software launches said CPU into execution of said operating system software, and so that a user can command said CPU, through said operating system software, to begin execution of an application software program;

system feature extension software, stored in said non-volatile memory, which contains a plurality of machine-dependent routines, and a self-describing feature table which contains pointers to said machine-dependent routines; and

multiple device driver programs, each accessible by said application software programs running on said CPU to define a software interface to specific features of at least one said I/O device; and wherein said device driver programs can make calls, to said machine-dependent routines, which are dependent on data in said self-describing feature table; and wherein said CPU, under control of said application software, programmably calls on one of said operating system a software, for interfacing to a selected one of said plural I/O devices according to a format which is substantially independent of a hardware type of said computer system, said CPU, under control of said operating system software, calling on said system feature extension software to provide a low-level interface to electrical operations; and

said device driver programs for interfacing to said selected I/O device according to a format which is substantially independent of said hardware type of said computer system;

wherein said systems of said first plurality are all mutually similar to each other, and said systems of said second plurality are all mutually similar to each other, and said systems of said third plurality are all mutually similar to each other;

wherein said systems of said first plurality each differ, in at least one hardware element, from every system of said second plurality;

wherein said systems of said first plurality each differ, in at least one hardware element, from every system of said third plurality;

wherein said systems of said second plurality each differ, in at least one hardware element, from every system of said third plurality;

wherein at least one of said device driver programs are able to make calls, to said machine-dependent routines, which are conditioned on data in said self-described feature table, and exists in the same form in all said computer systems, of said first, second, and third pluralities.

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